

Graph theory

Flow problems

Flows: definitions and properties 1

- Flow networks are a fundamental concept in graph theory used to model the movement of resources through interconnected nodes.
- A flow network is a directed graph where a certain homogeneous quantity can stream between connected vertices through edges. This quantity is called flow.
- Each edge is characterized by a capacity, a non-negative real number.
- There are also two special vertices:
- $s \in V$ is the **source** vertex from which the flow originates.
- $t \in V$ is the **sink** vertex where the flow is collected.
 - No arc reaches the source.
 - No arc leaves the sink.

Flows: definitions and properties 2

- A flow non-negative real-valued function f defined on the arcs satisfying:
 - *Capacity constraint* : $f(a) \leq c(a)$;
 - *Conservation of flow* : for any vertex other than s and t , the sum of the flows on the incoming edges and the sum of the flows on the outgoing edges are equal.
 - $f(x,y) = -f(y,x)$
- Examples: electrical or hydraulic circuits, communication networks, transport modeling

Flows: definitions and properties 3

- the sum of the flows on the arcs leaving the source and the sum of the flows on the arcs arriving at the sink are equal;
 - This value is the value of the flow $|f|$;
- if we separate the vertices into two subsets **E** containing **s** and **F = A - E**, containing **t**, then the sum of the values of the flow on the arcs from **E** to **F** minus the sum of the values of the flow on the arcs from **F** to **E** is also $|f|$.
- Such a separation into two subsets of the vertices is called a cut and this difference in flow sums is called the net flow crossing the cut.

Ford and Fulkerson Algorithm

Max_Flow (G)

```
{  
    MF=0; //the value of the max flow  
    While (there exists an augmenting path in G) Do  
    {  
        Find a path increasing P  
         $Cf(P)$  = smallest capacity on P  
        MF=MF+  $Cf(P)$   
        For each arc ( u,v ) in P {  
             $cf(u,v) = cf(u,v) - Cf(P)$   
             $cf(seen) = cf(seen) + Cf(P)$   
        }  
    }  
}
```

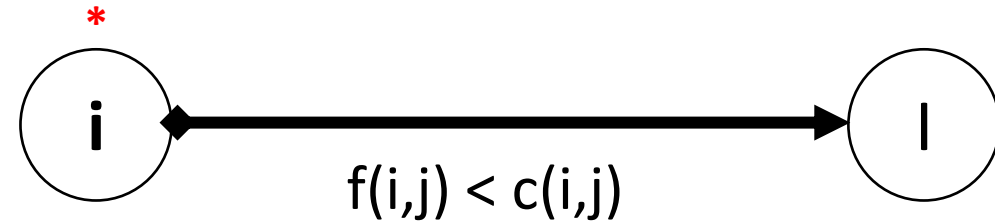
Definition: increasing path

- A path C between “ s ” and “ t ” is said to be increasing with respect to a flow $f = (f(i, j), (i, j) \in E)$ feasible between $\underline{s}$ and $\underline{t}$ if:
 - $f(i, j) < c(i, j)$ if $(i, j) \in C^+ ((i, j) \in E)$ conformal arc
 - $f(i, j) > 0$ if $(i, j) \in C^- ((j, i) \in E)$ non-conforming arc
- C^+ is the set of all arcs of C that meet in the right direction.
- C^- is the set of arcs of C that meet in the opposite direction.

Marking procedure

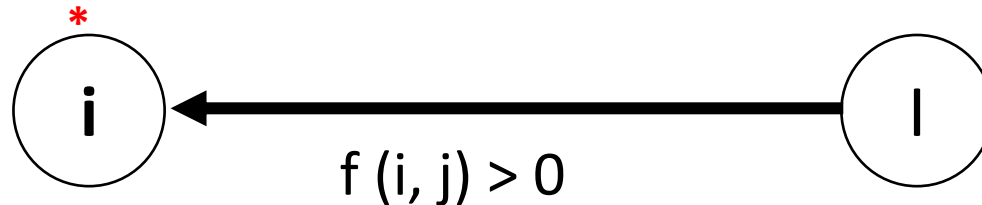
- Direct marking:

- If for an arc (i, j) we have: i marked, j unmarked and $f(i, j) < c(i, j)$ Then we mark j by $(i, +)$ and we set $\delta(j) = \min(\delta(i), c(i, j) - f(i, j))$

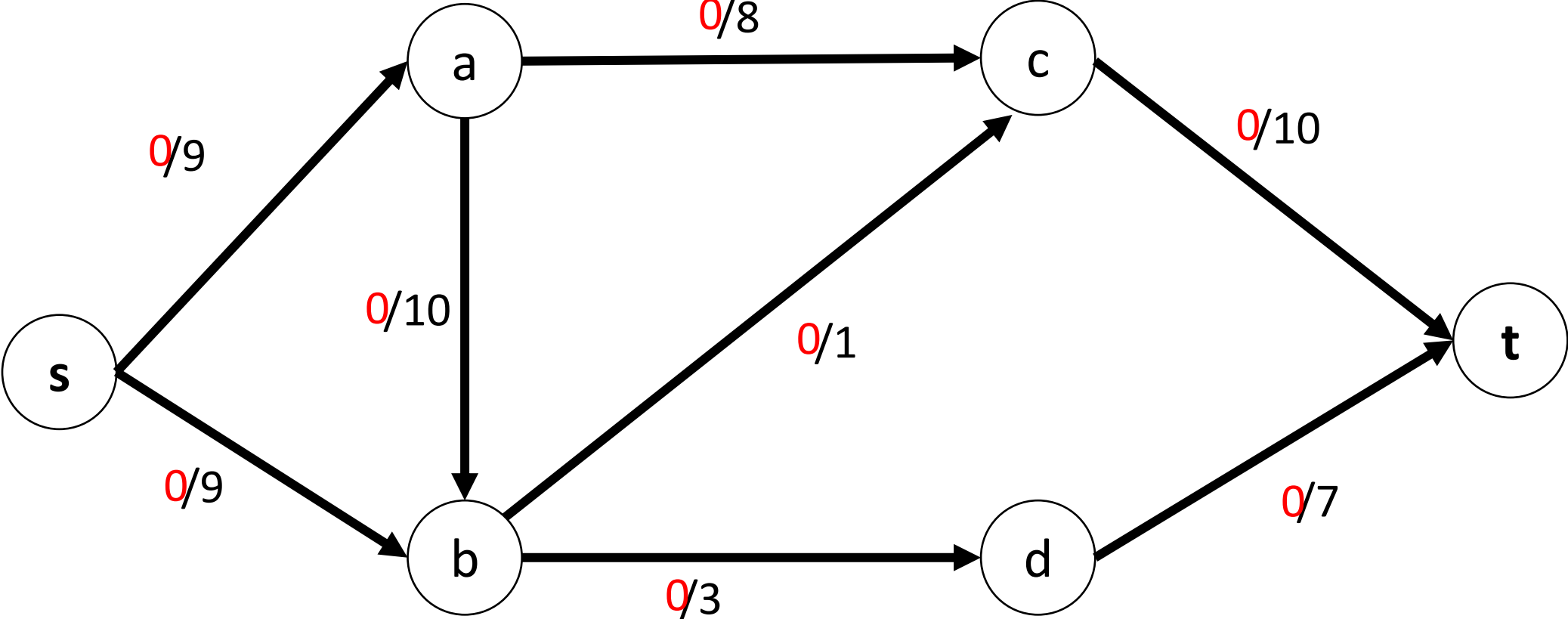


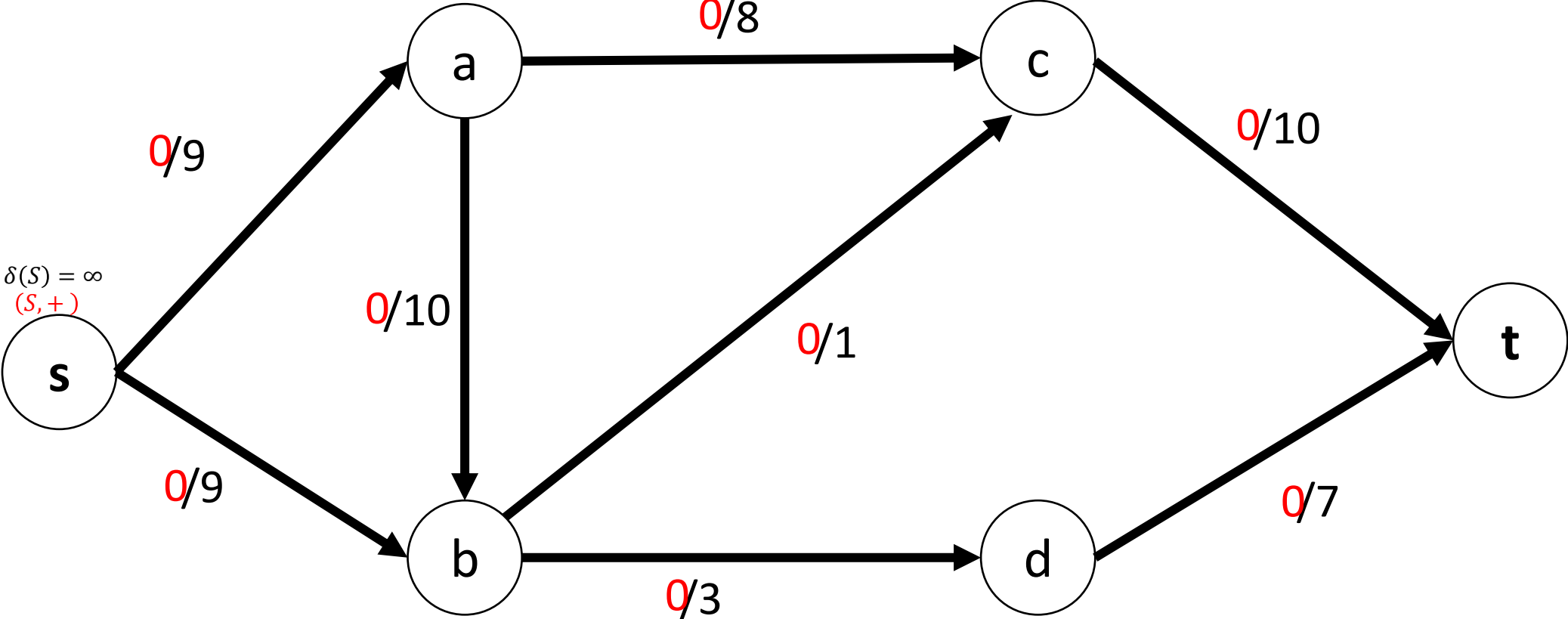
- Indirect marking:

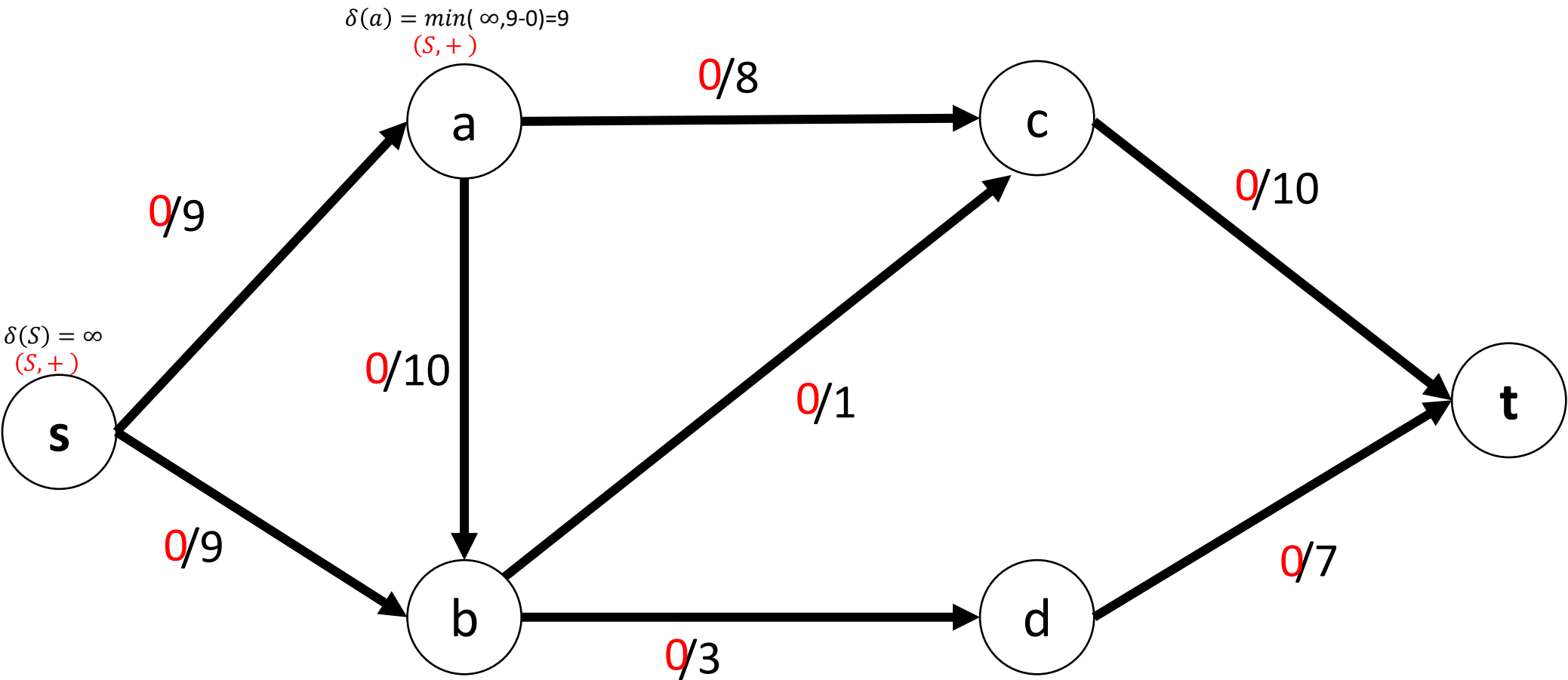
- If for an arc (j, i) we have: i marked, j unmarked and $f(j, i) > 0$ Then we mark j by $(i, -)$ and we set $\delta(j) = \min(\delta(i), f(j, i))$

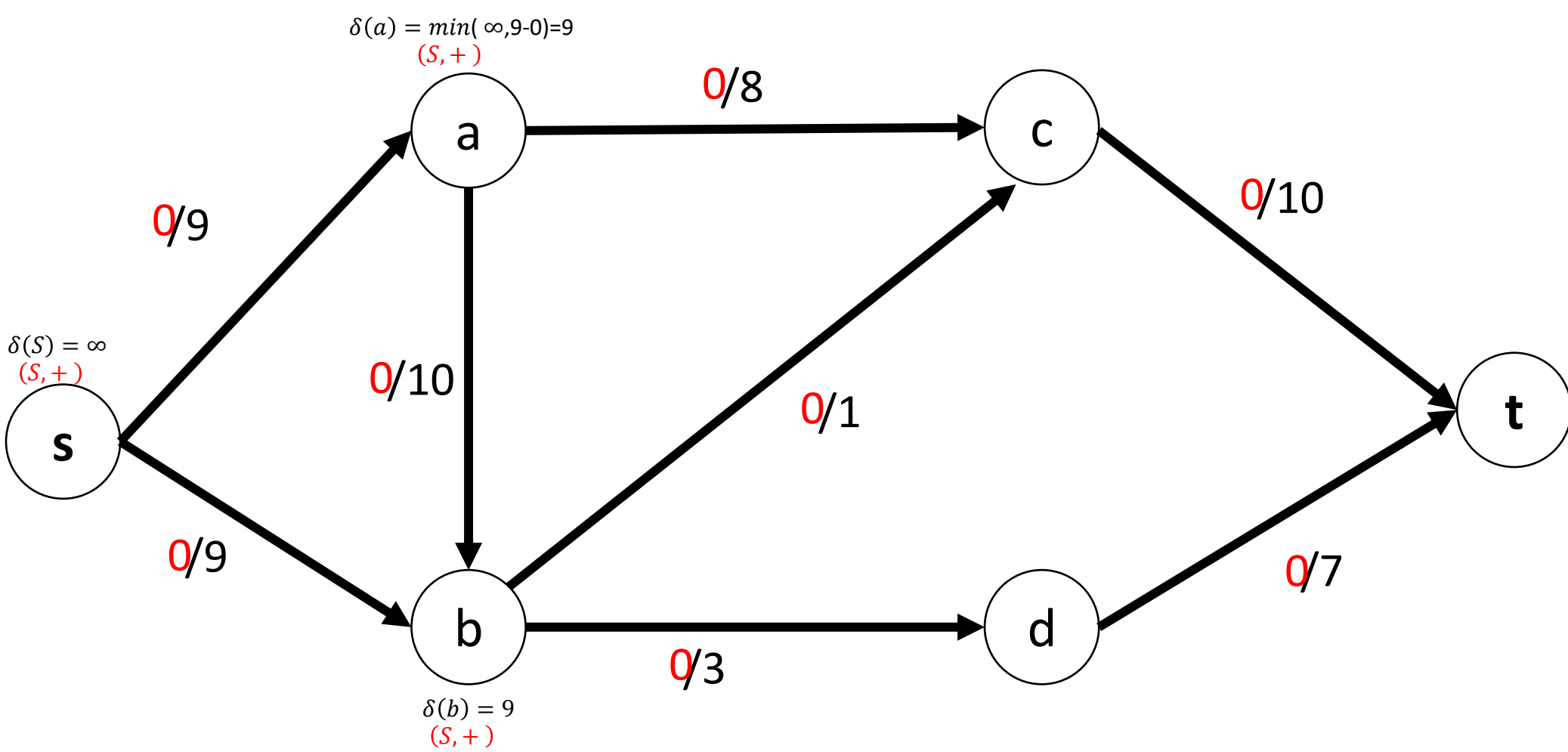


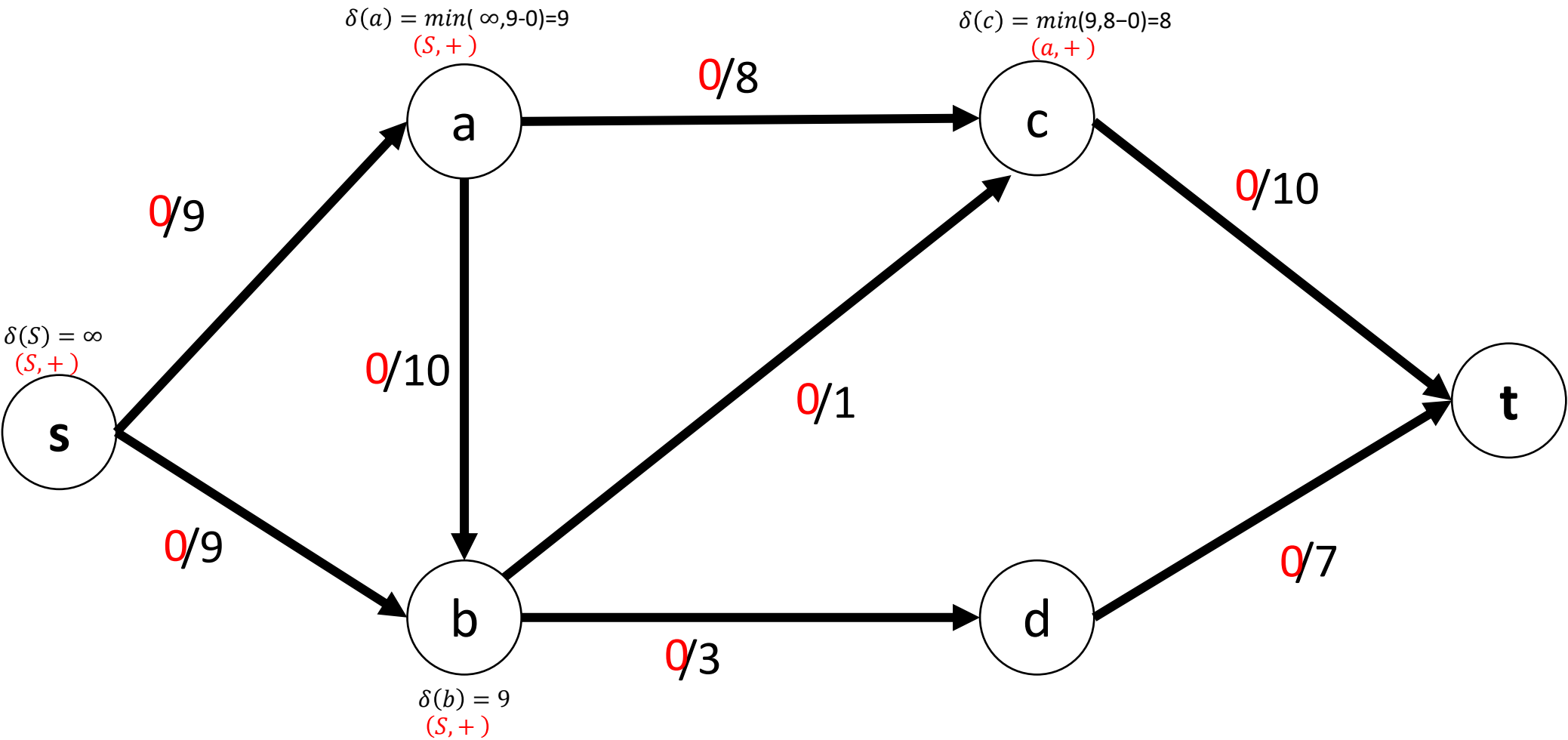
- $\delta(j)$ is the maximum flow you can increase from s to j .
- $\delta(i)$ is a value associated with vertex i , initialized to infinity for s .

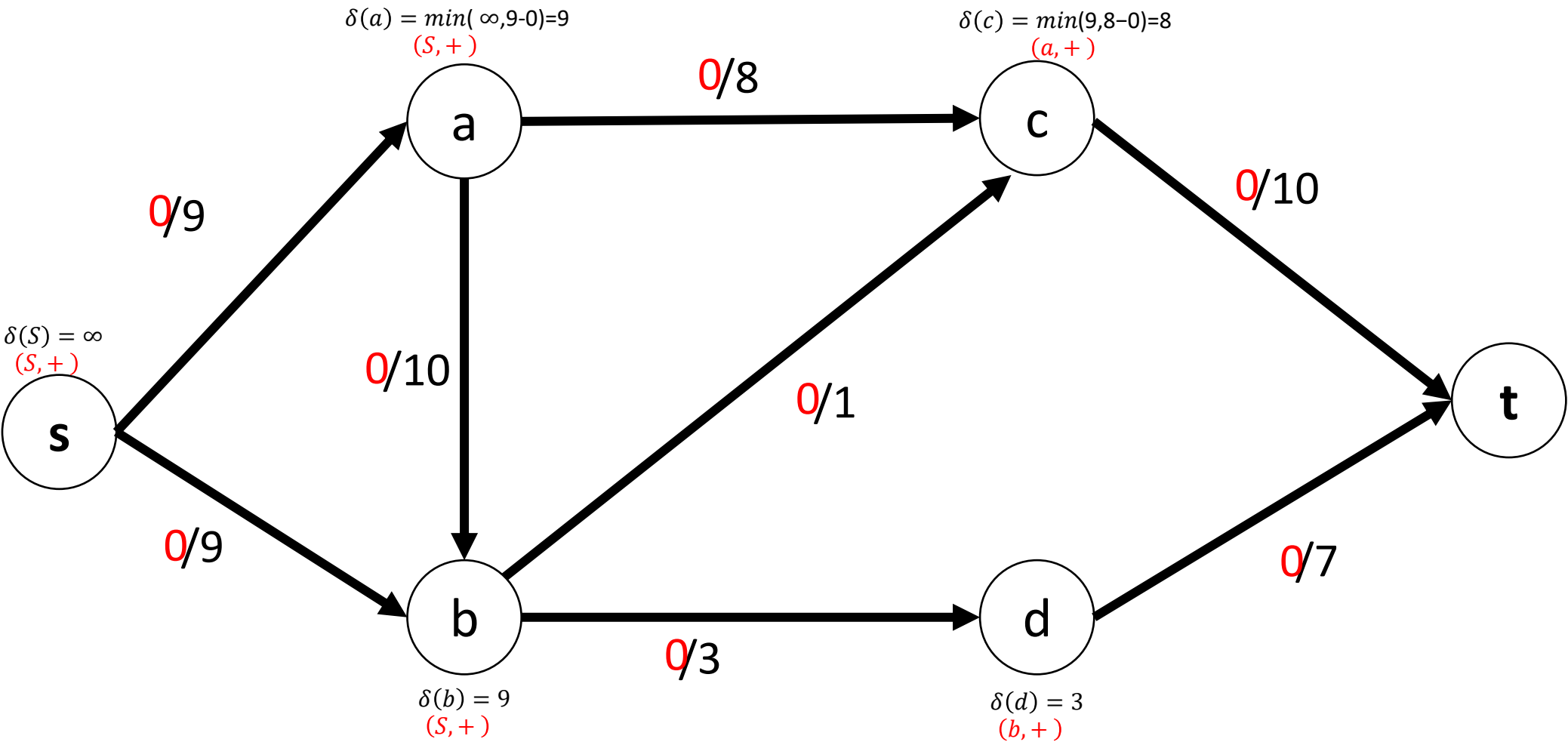


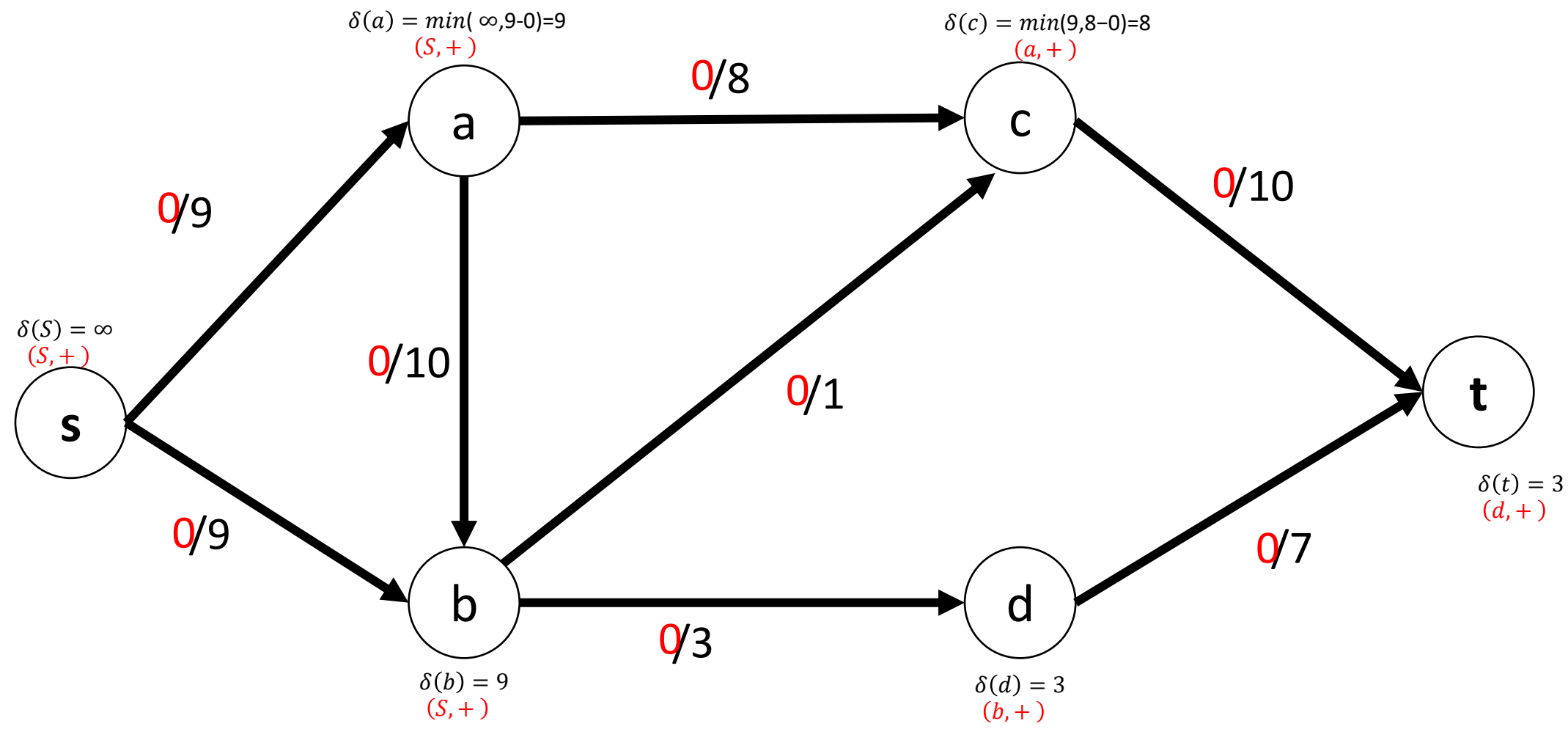


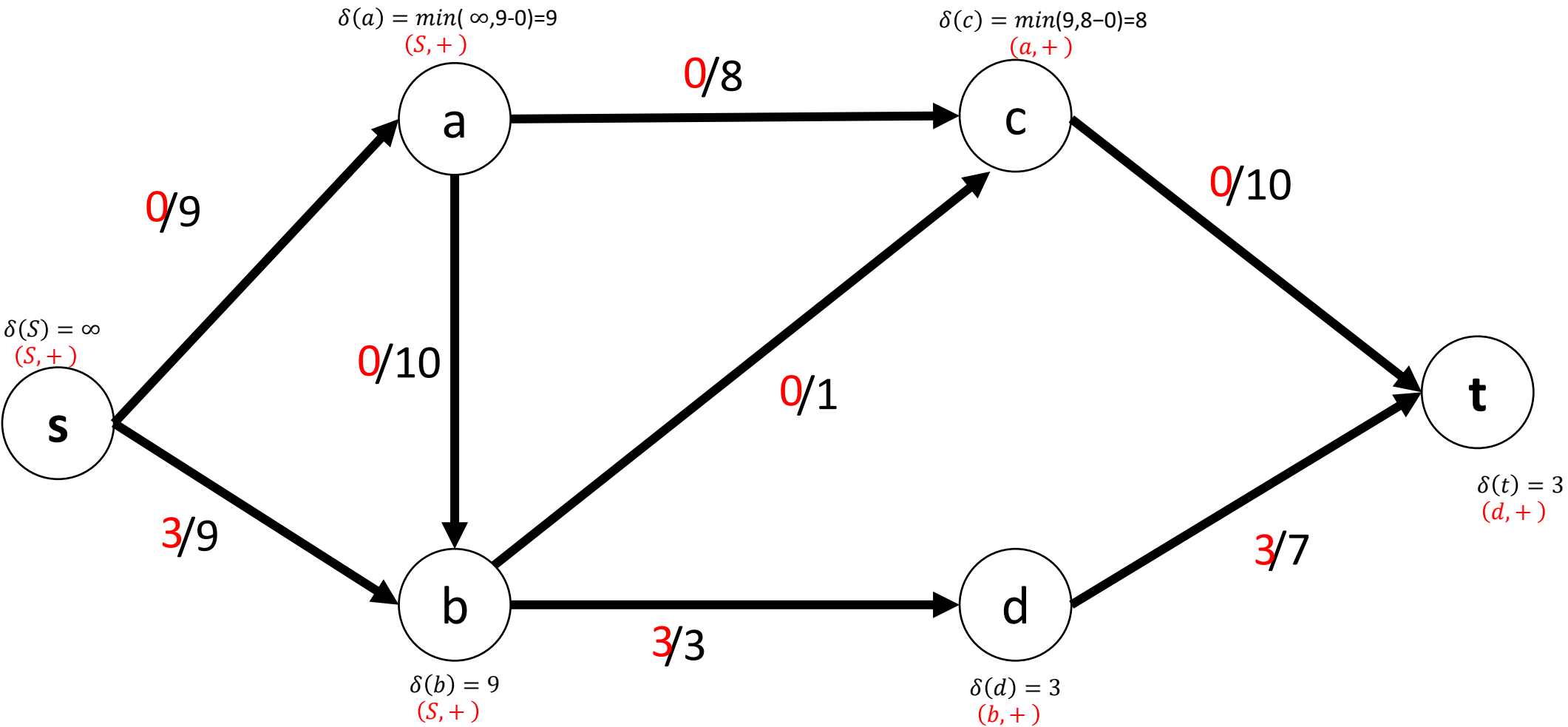


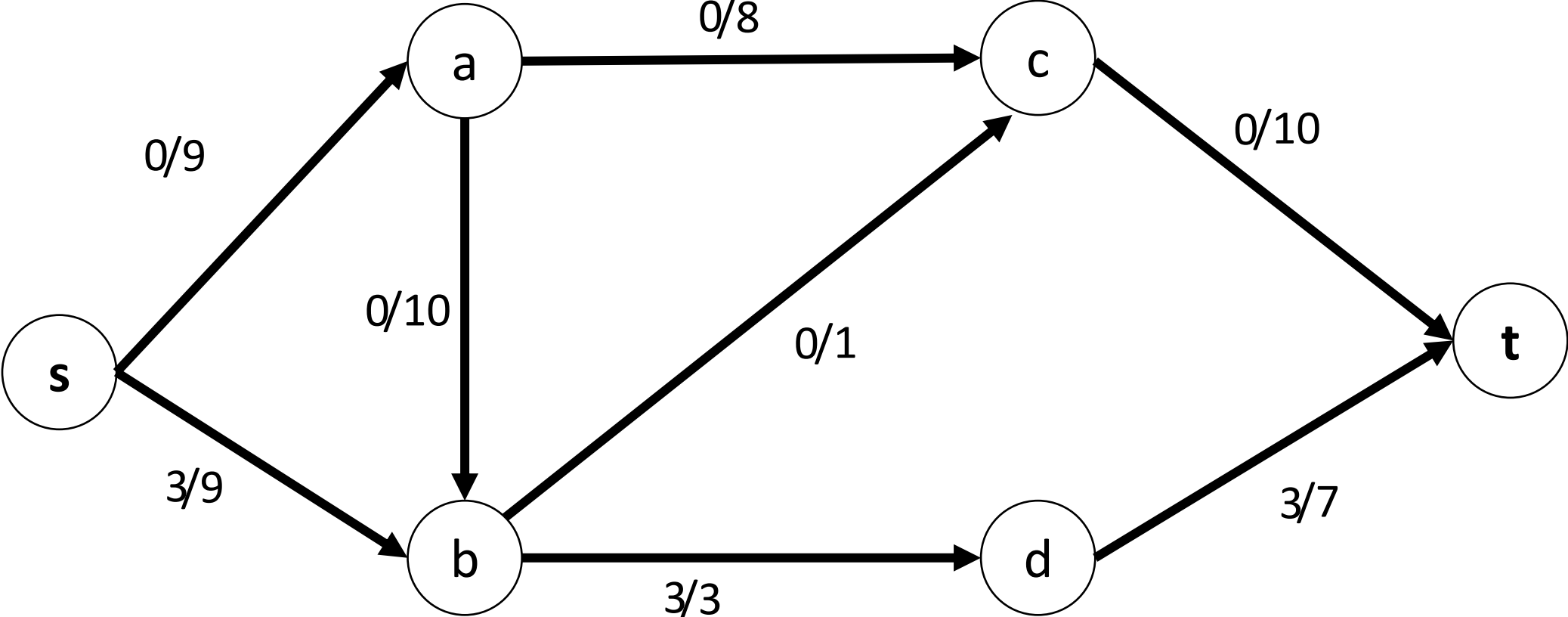


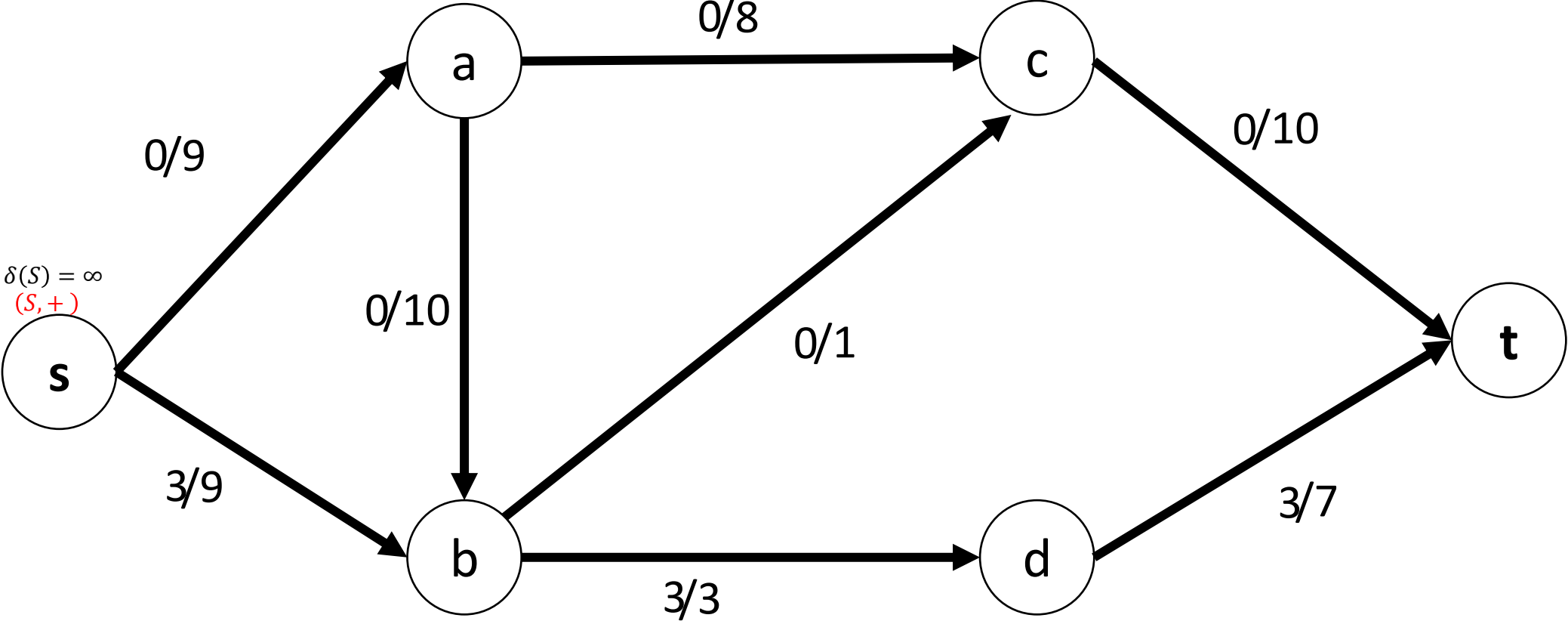


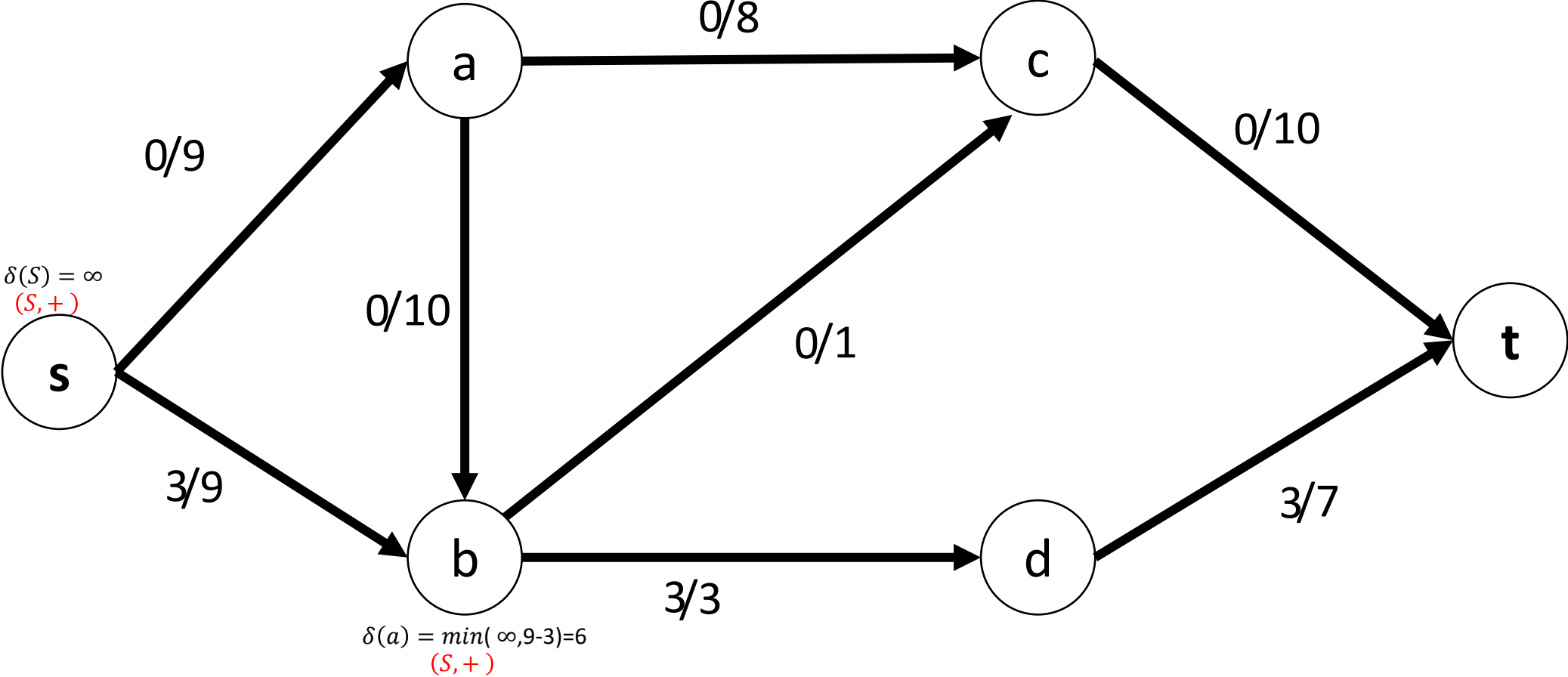


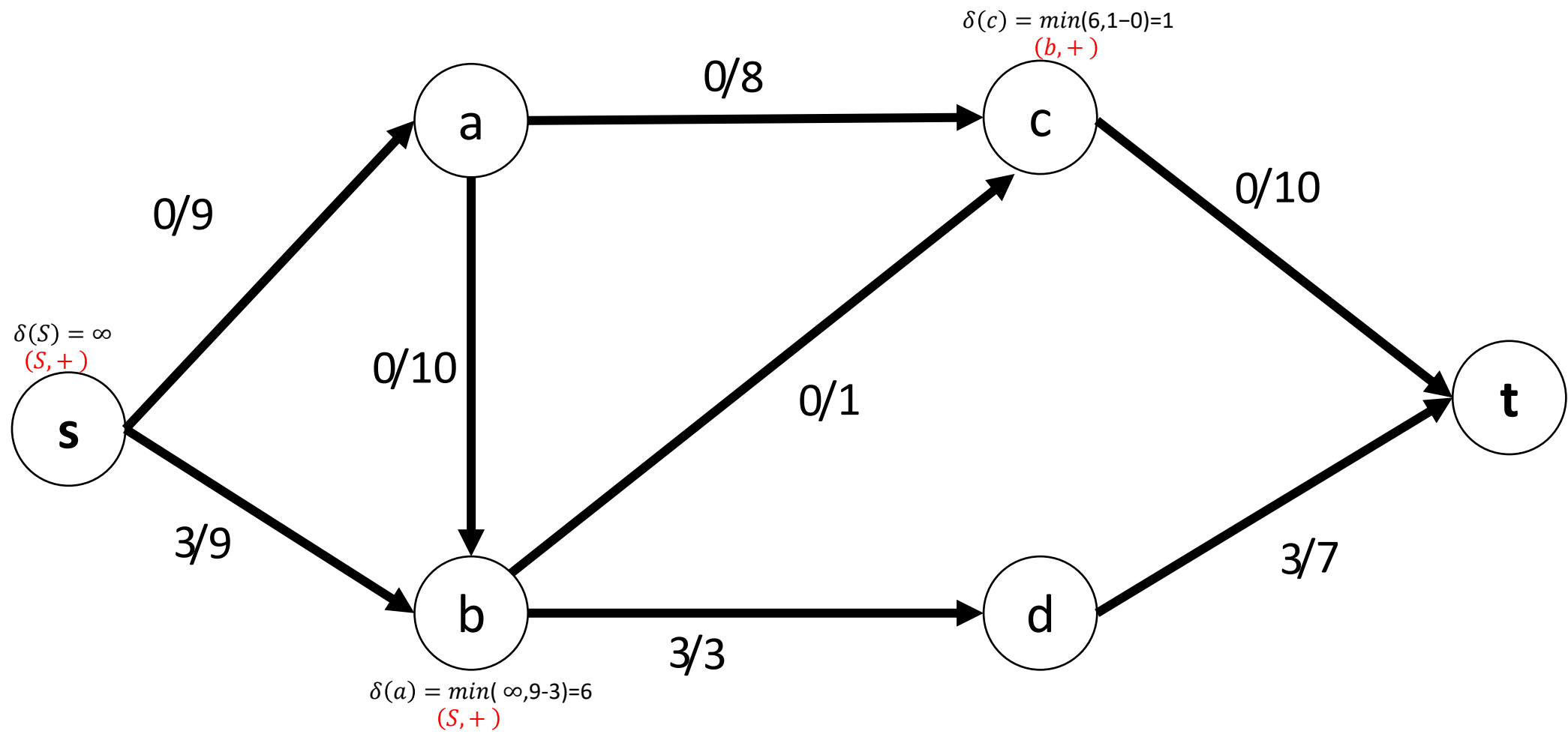


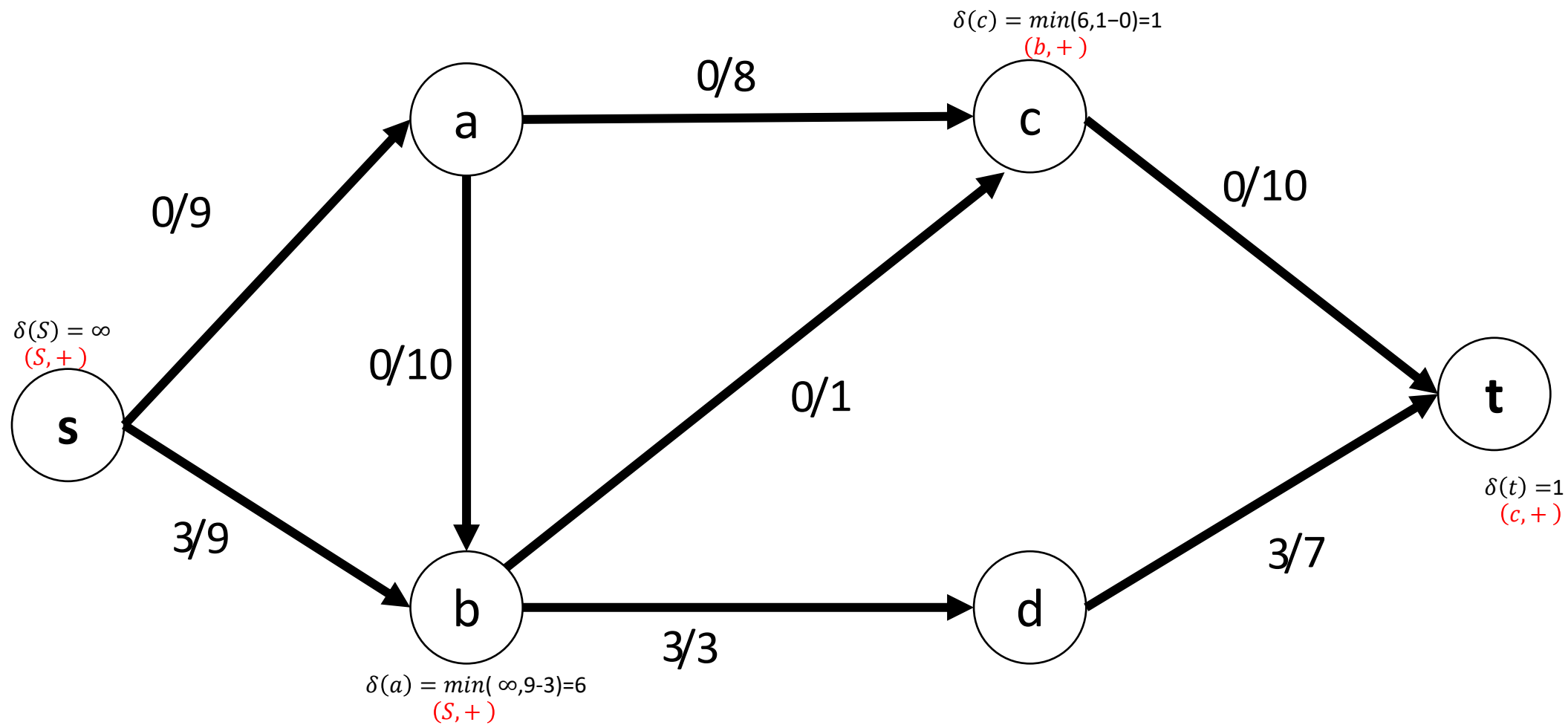


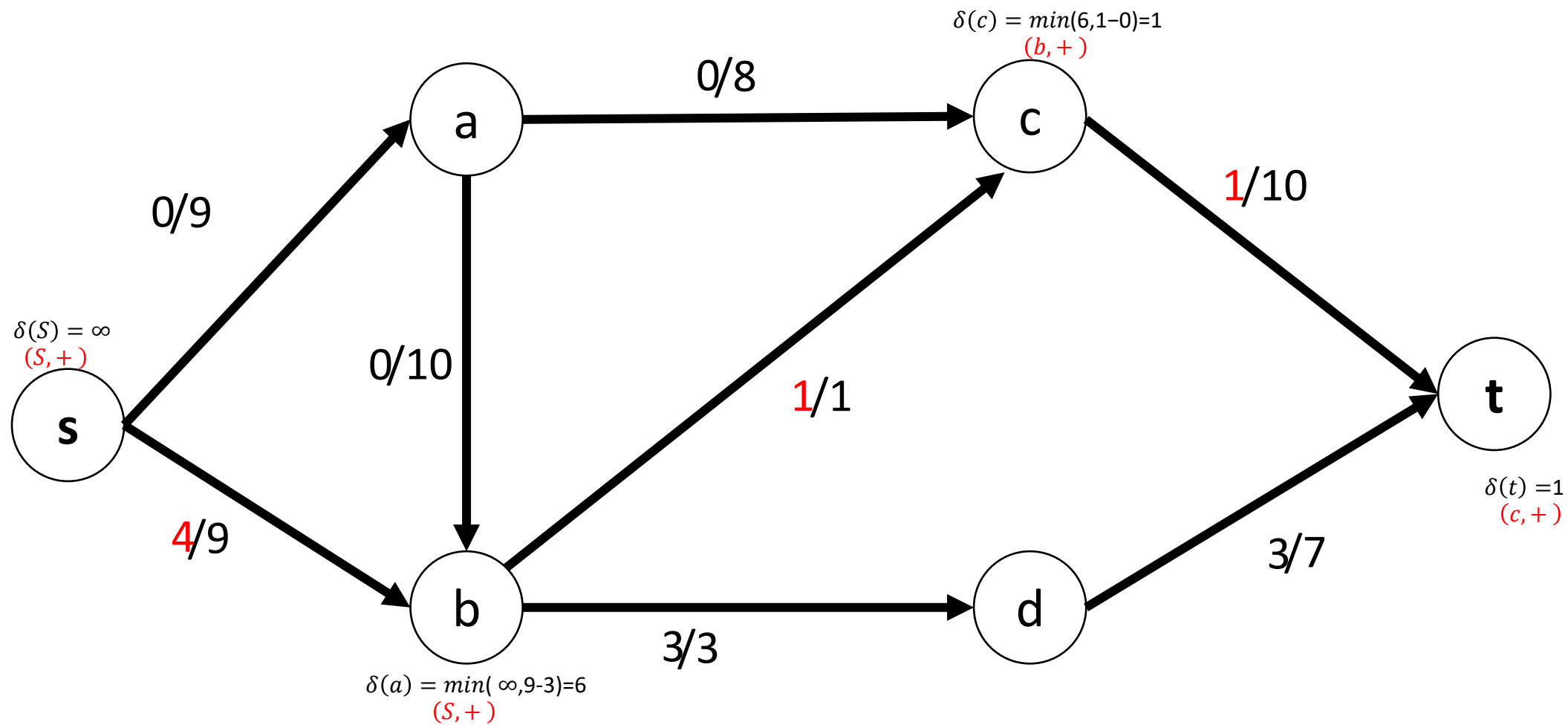


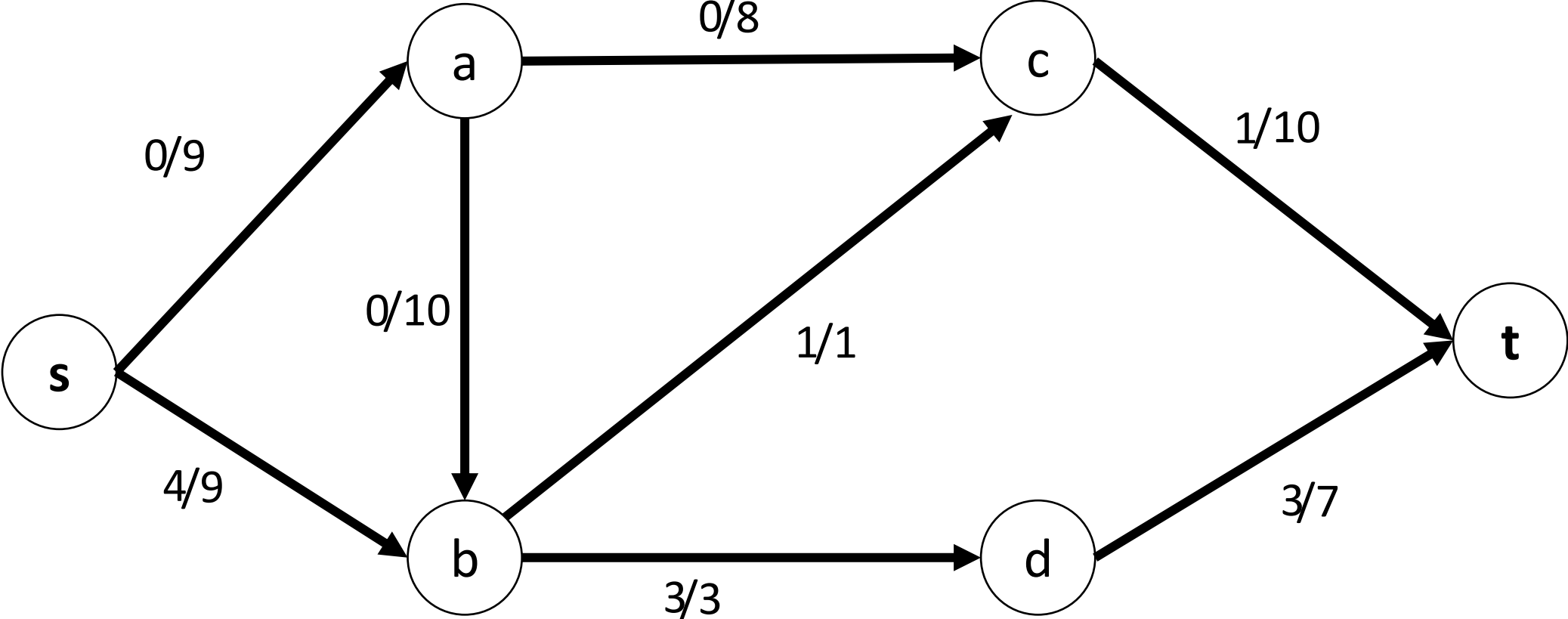


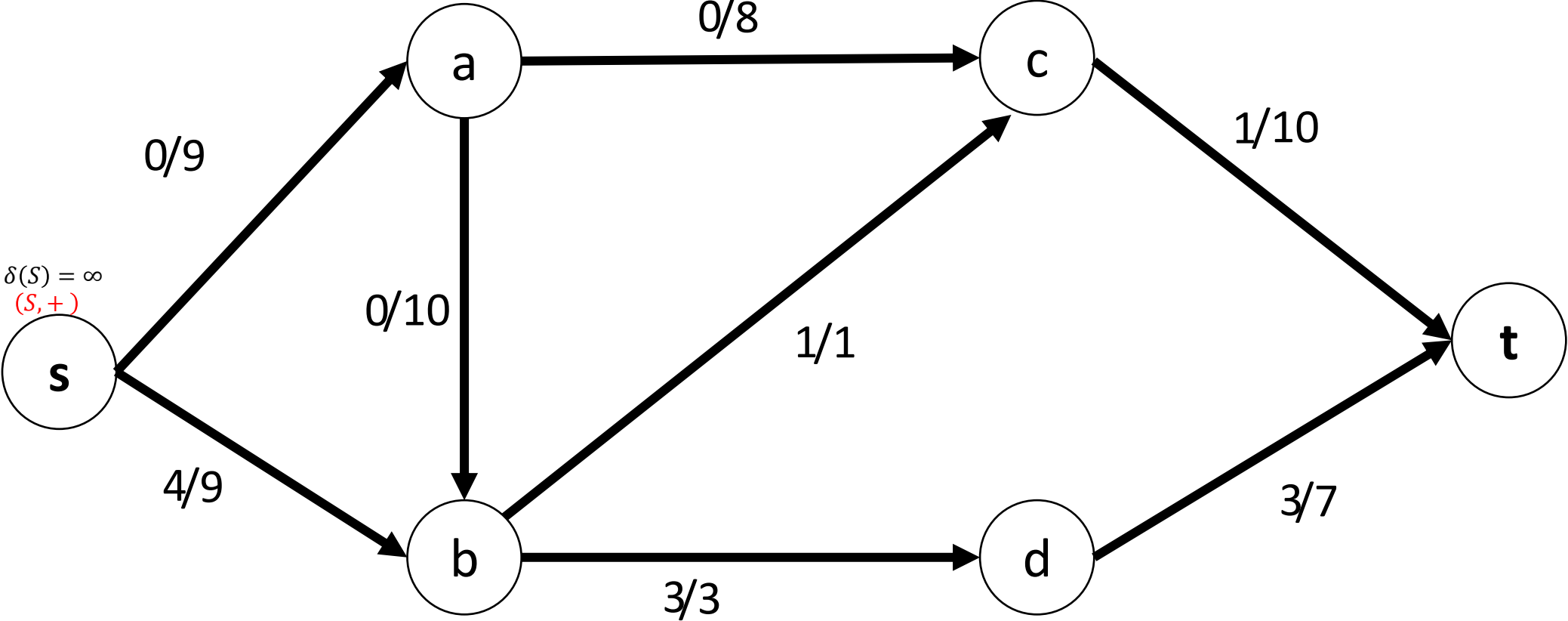


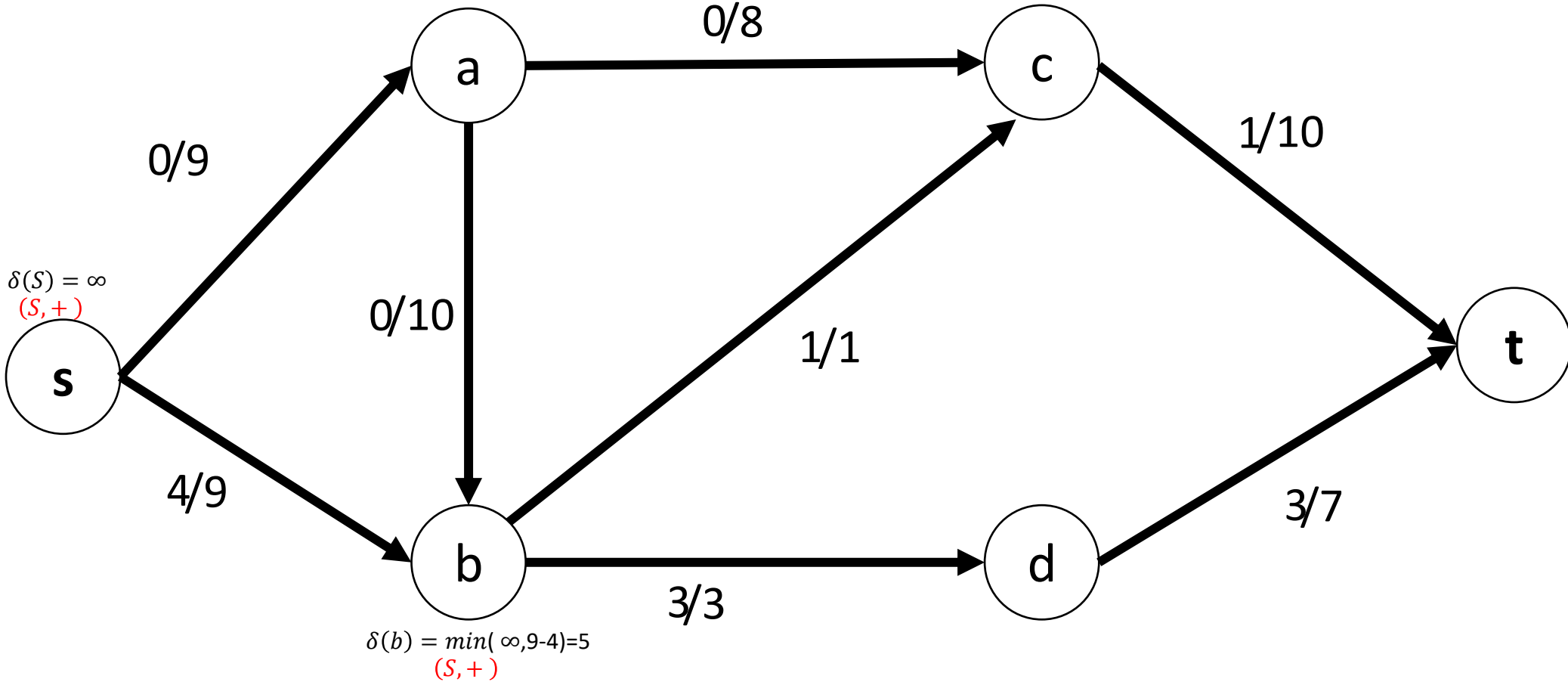


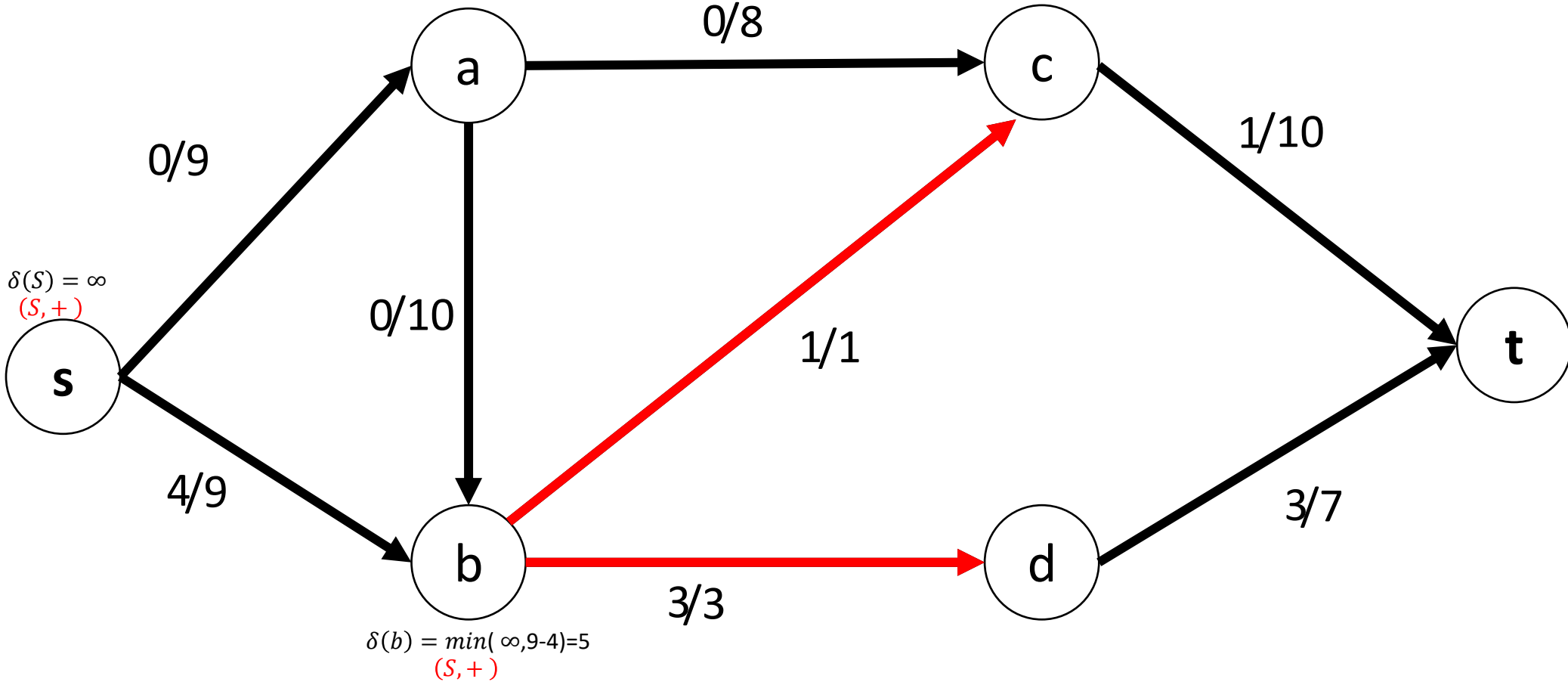


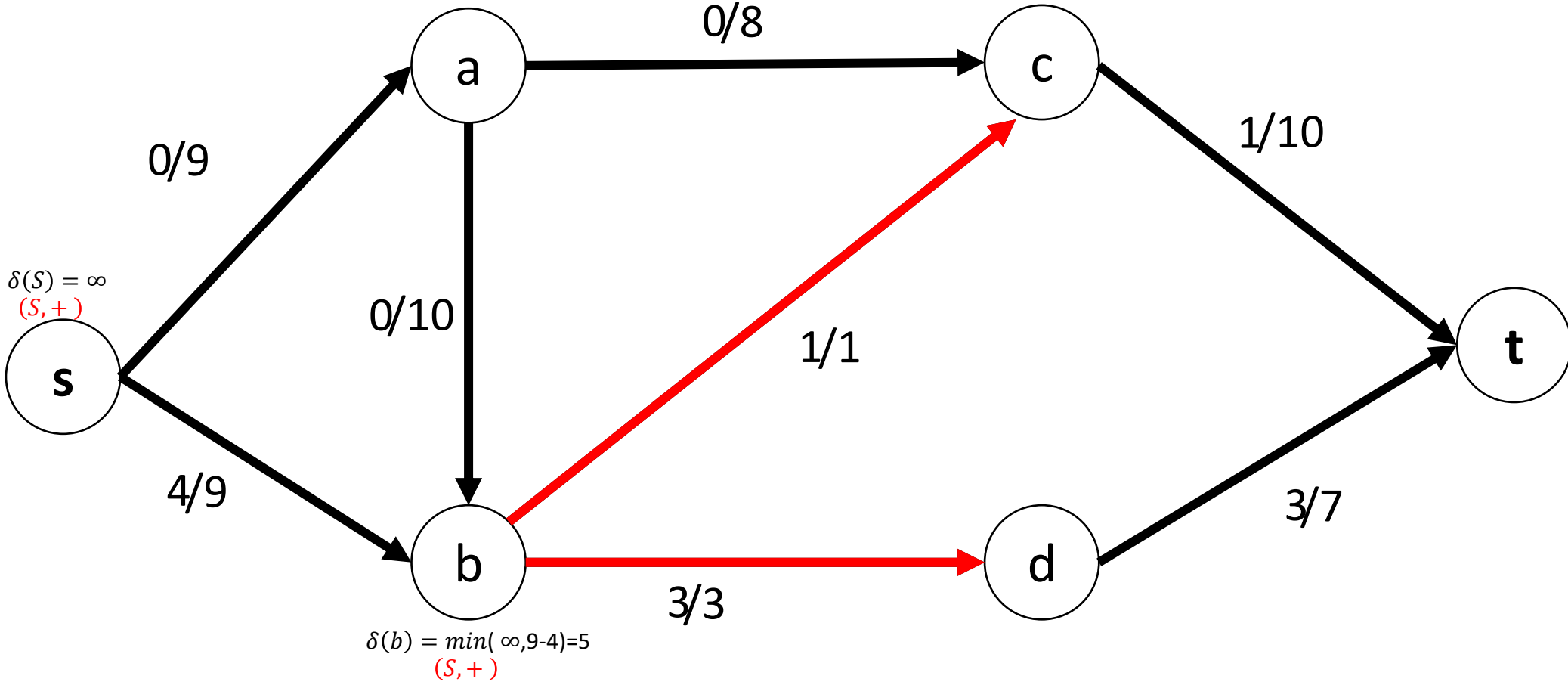


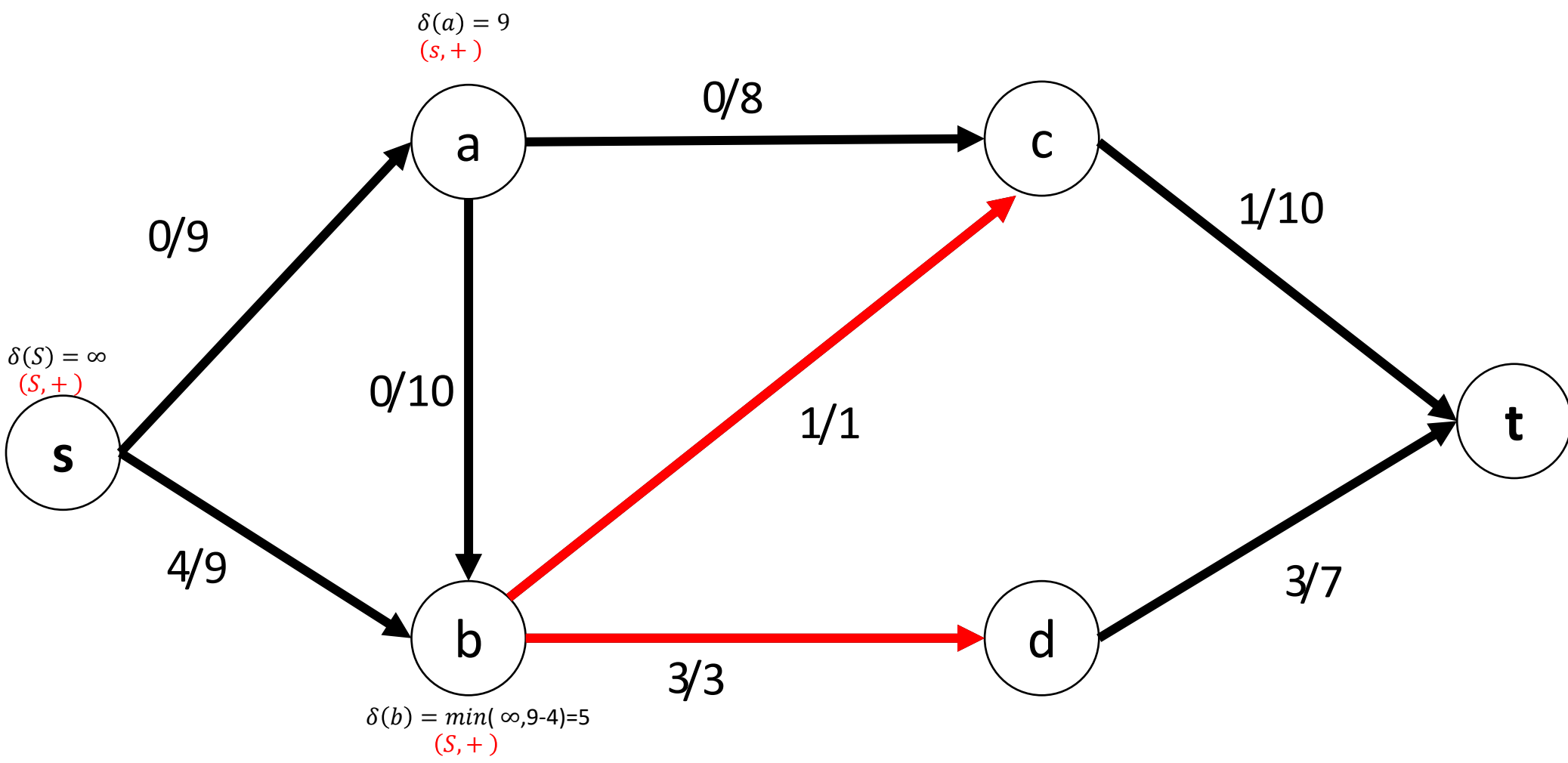


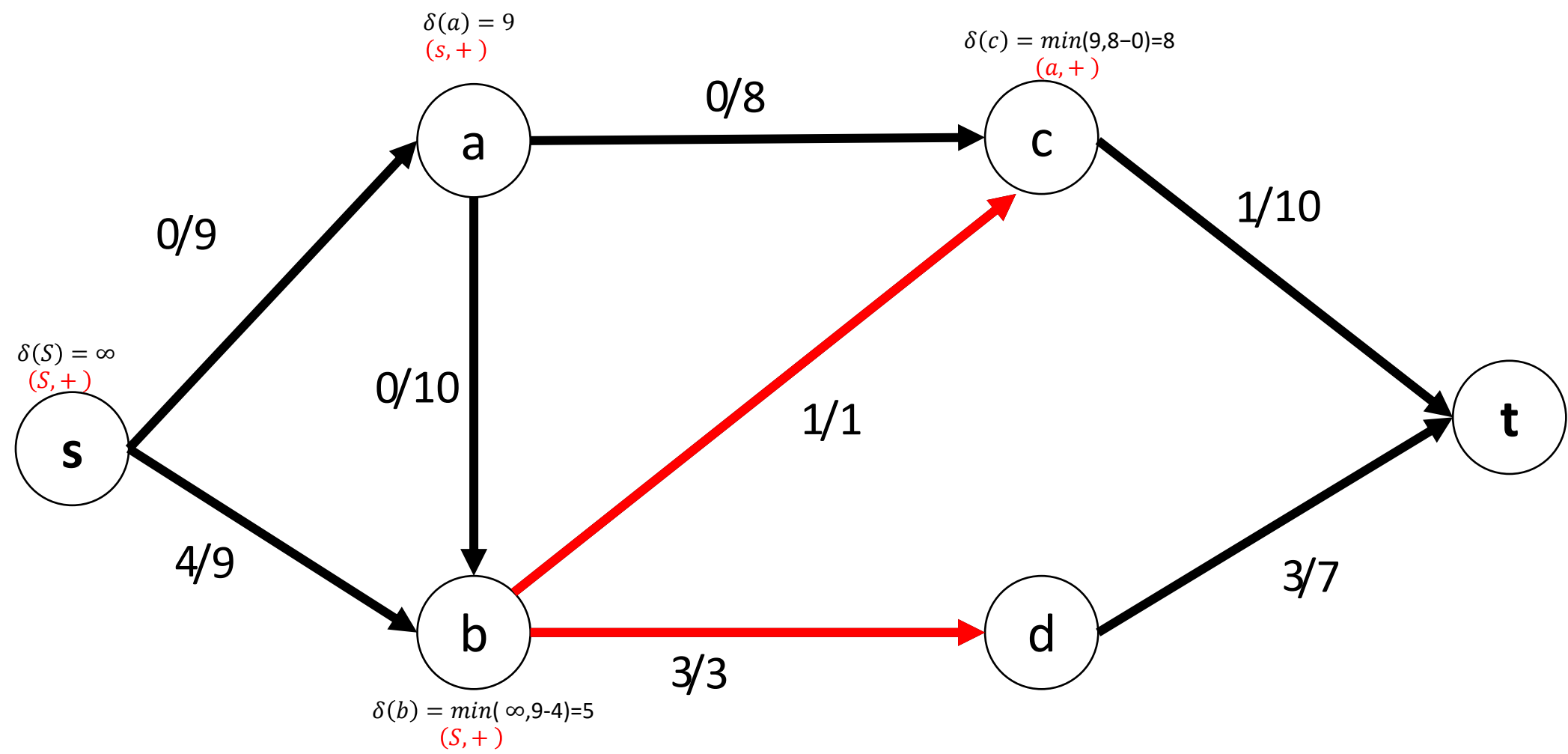


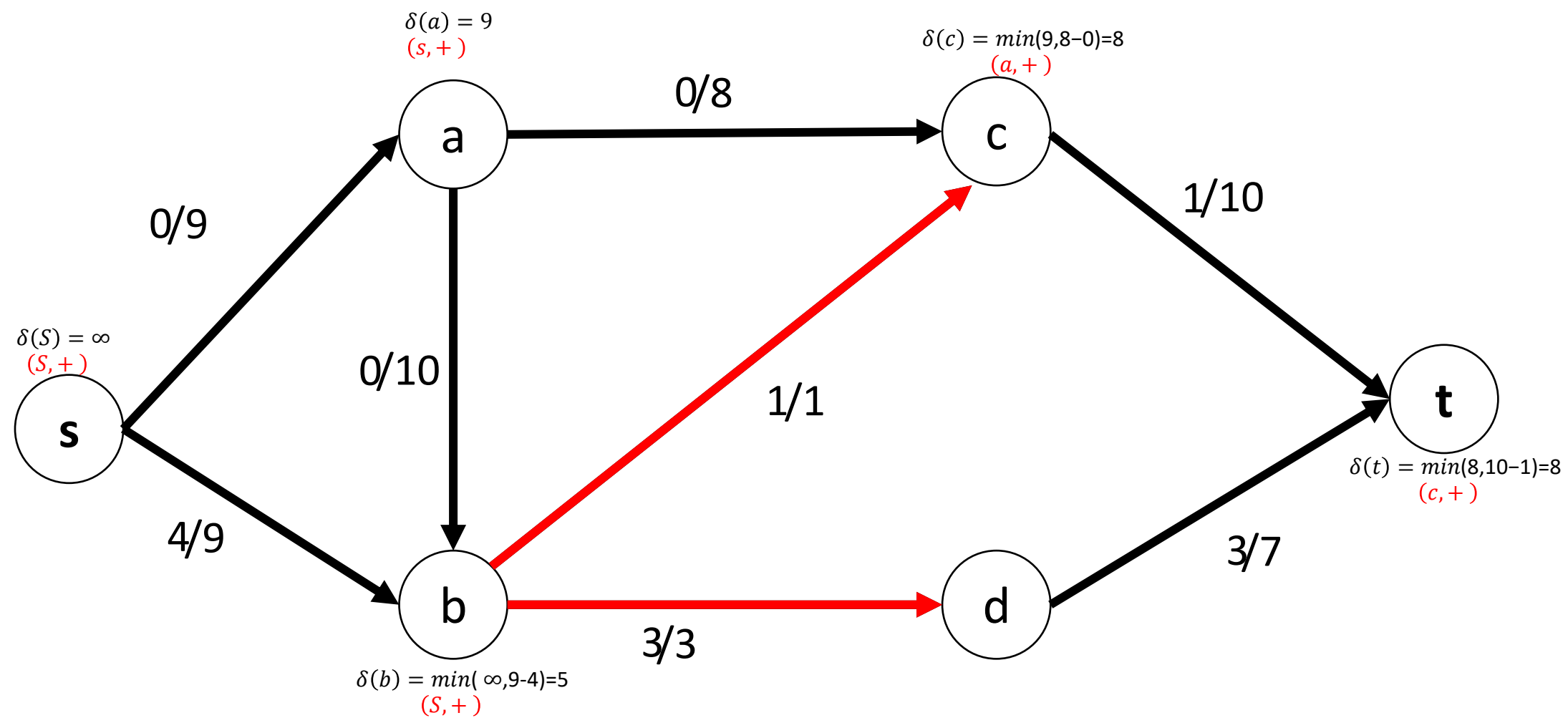


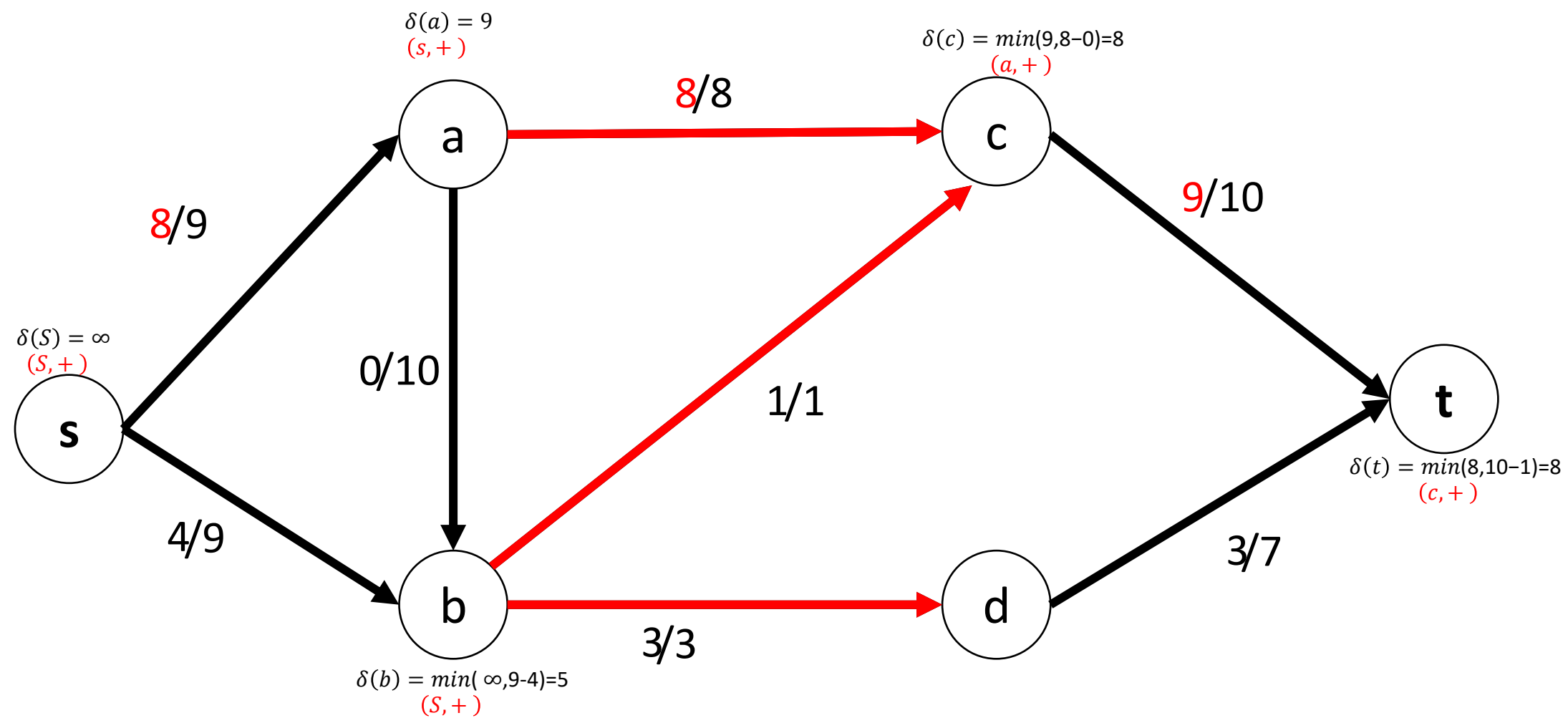












$$C(f) = 9 + 3 = 12$$

Definitions

- **Cut** : A cut is a partition of the set of vertex V into two subsets S and $S'=V-S$. The cut is denoted $[S,S']$.
- A cut like the set of arcs whose ends are in the different partitions S and S' .
- A cut is called a **st cut** if $s \in S$ and $t \in S'$.
- An arc (i,j) such that $i \in S$ and $j \in S'$ is a *forward arc*, and
- An arc (i,j) such that $i \in S'$ and $j \in S$ is a *backward arc* of the cut $[S,S']$.
- The **capacity** of a cut **st cut** (S,S') is the sum of the capacities of the front arcs of this cut.

$$C(\delta(S)) = \sum_{(i,j) \in \delta(S)} C(i,j)$$

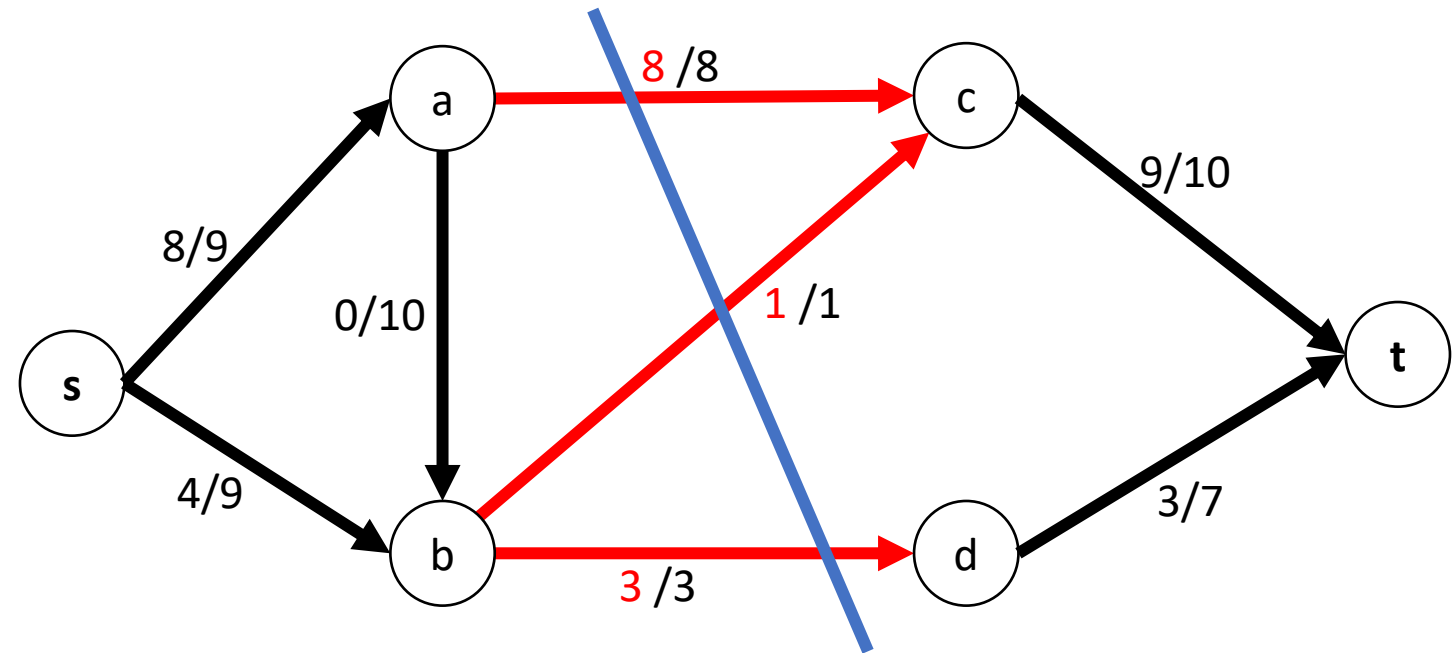
Minimal Cut

- A **st cut** whose capacity is minimum among all *st cuts* is a **minimum cut**.
- Any flow from **s** to **t** must pass from **S** to **S'** at some point, and thus uses the capacity of the edges from **S** to **S'**.
- This means that each cut of the graph puts a bound (limit) on the value of the maximum possible flow.

$$S = \{s, a, b\}, S' = \{c, d, t\}$$

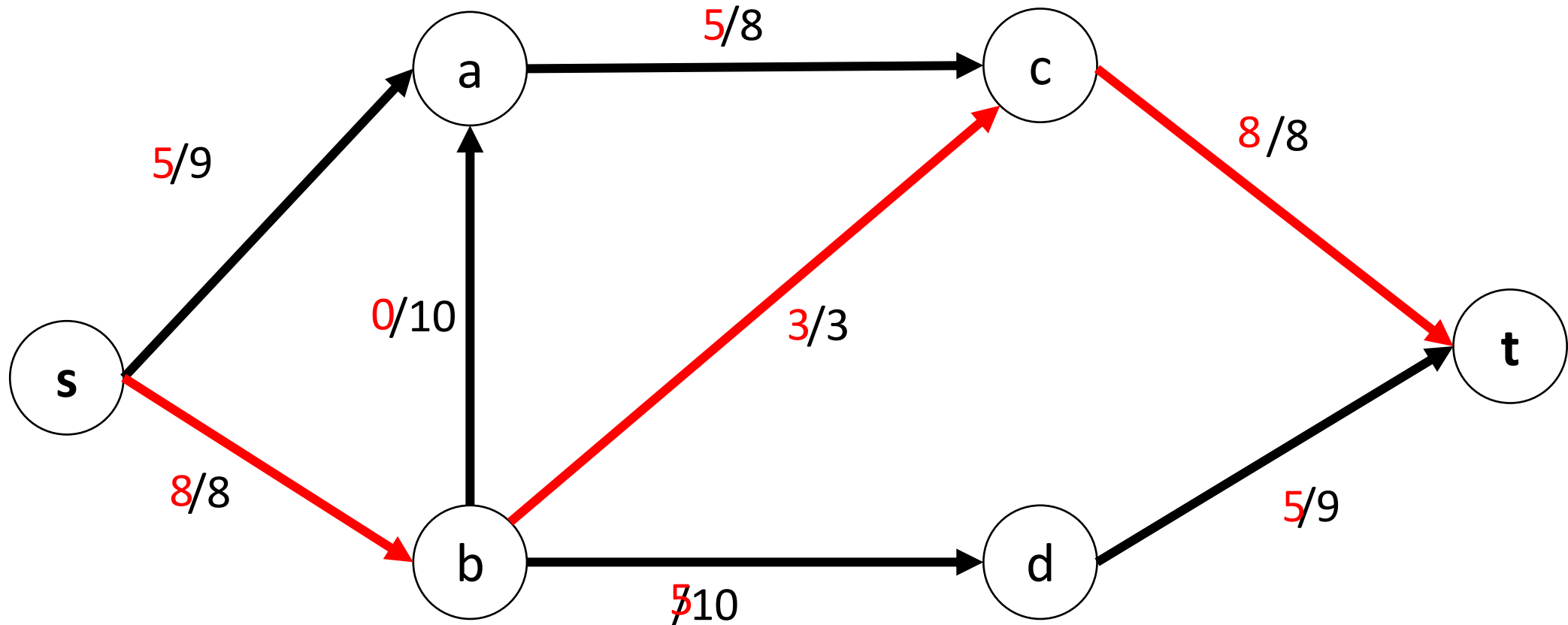
$$\delta(S) = \{(a, c), (b, c), (b, d)\}$$

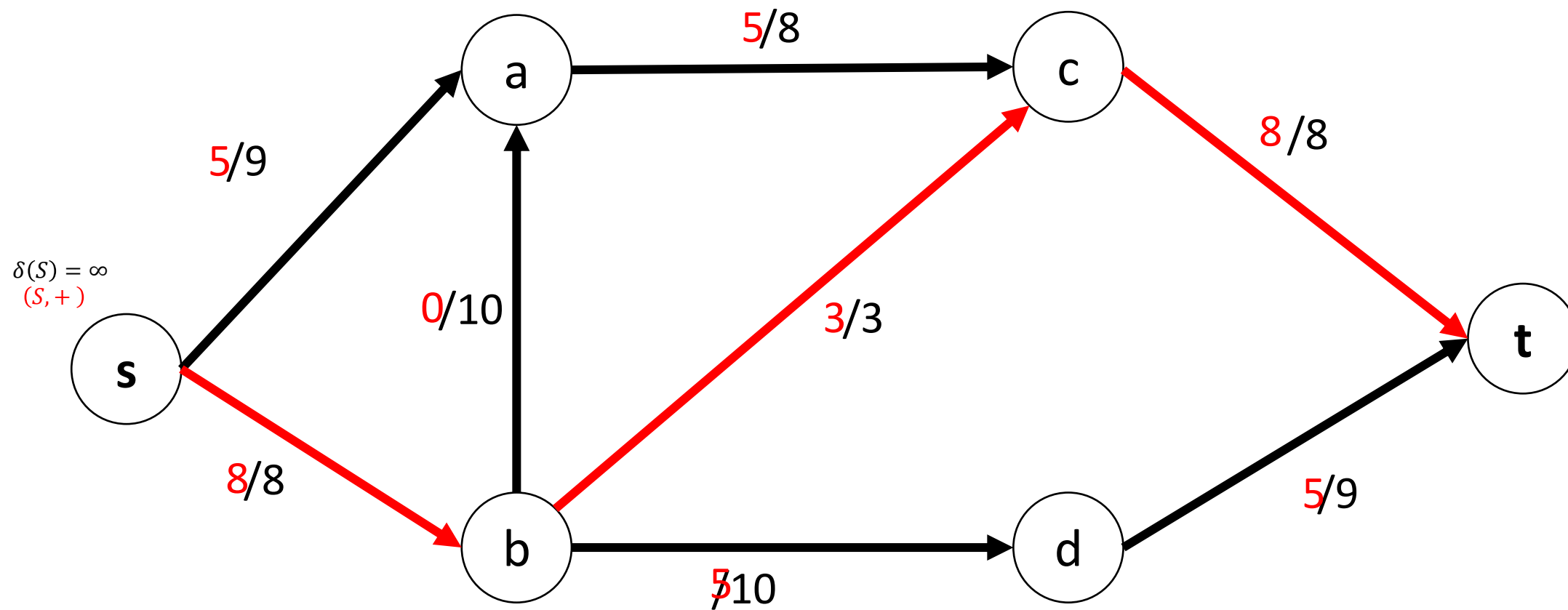
$$C(\delta(S)) = \sum_{(i,j) \in \delta(S)} C(i,j) = 8+1+3 = 12$$

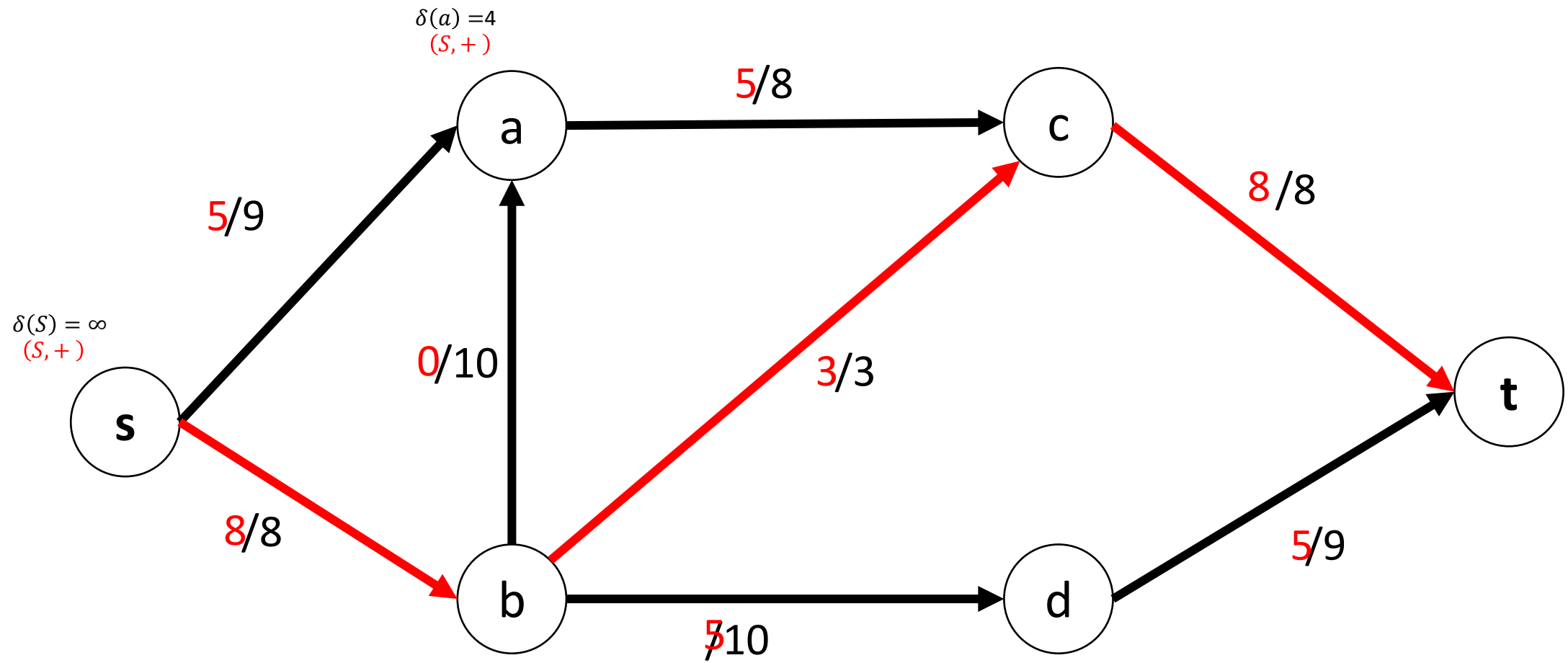


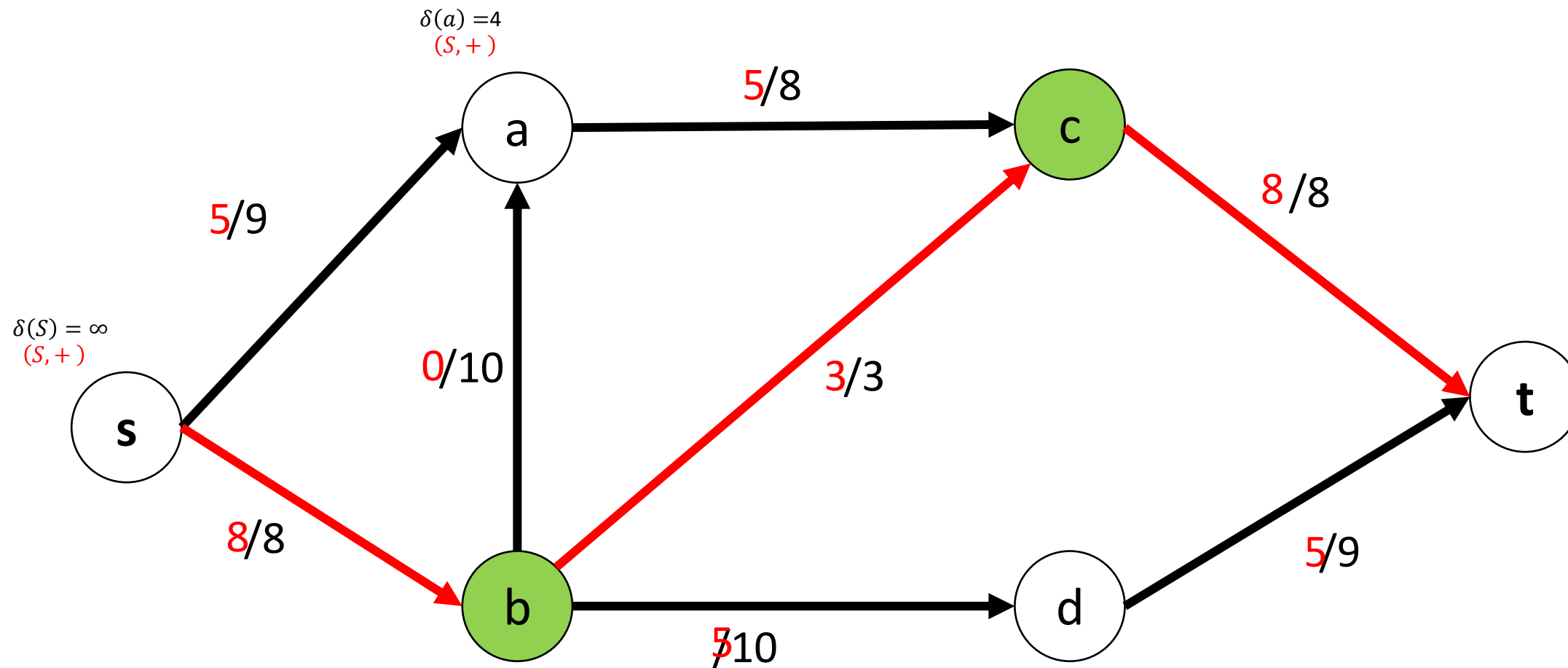
Exercises 1 :

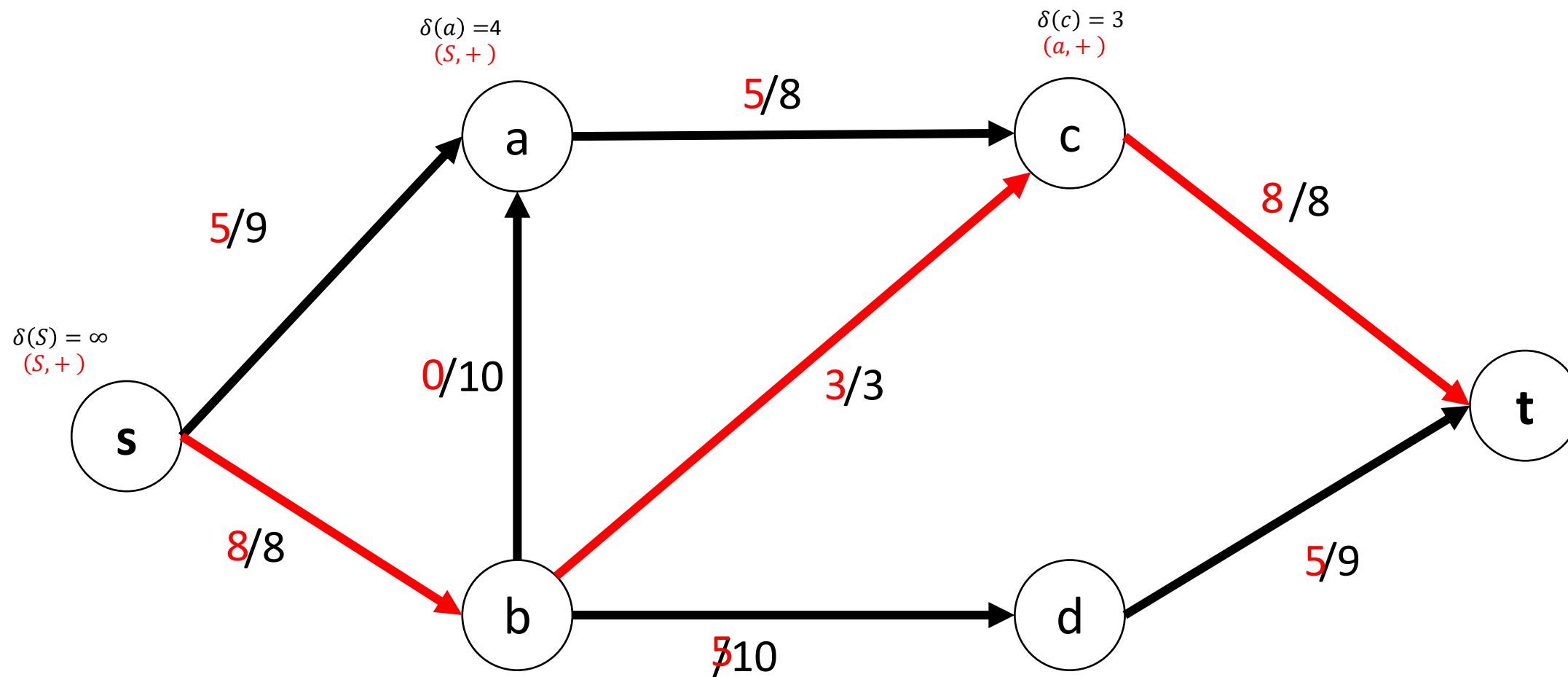
Find the maximum flow by completing the Ford-Fulkerson algorithm .

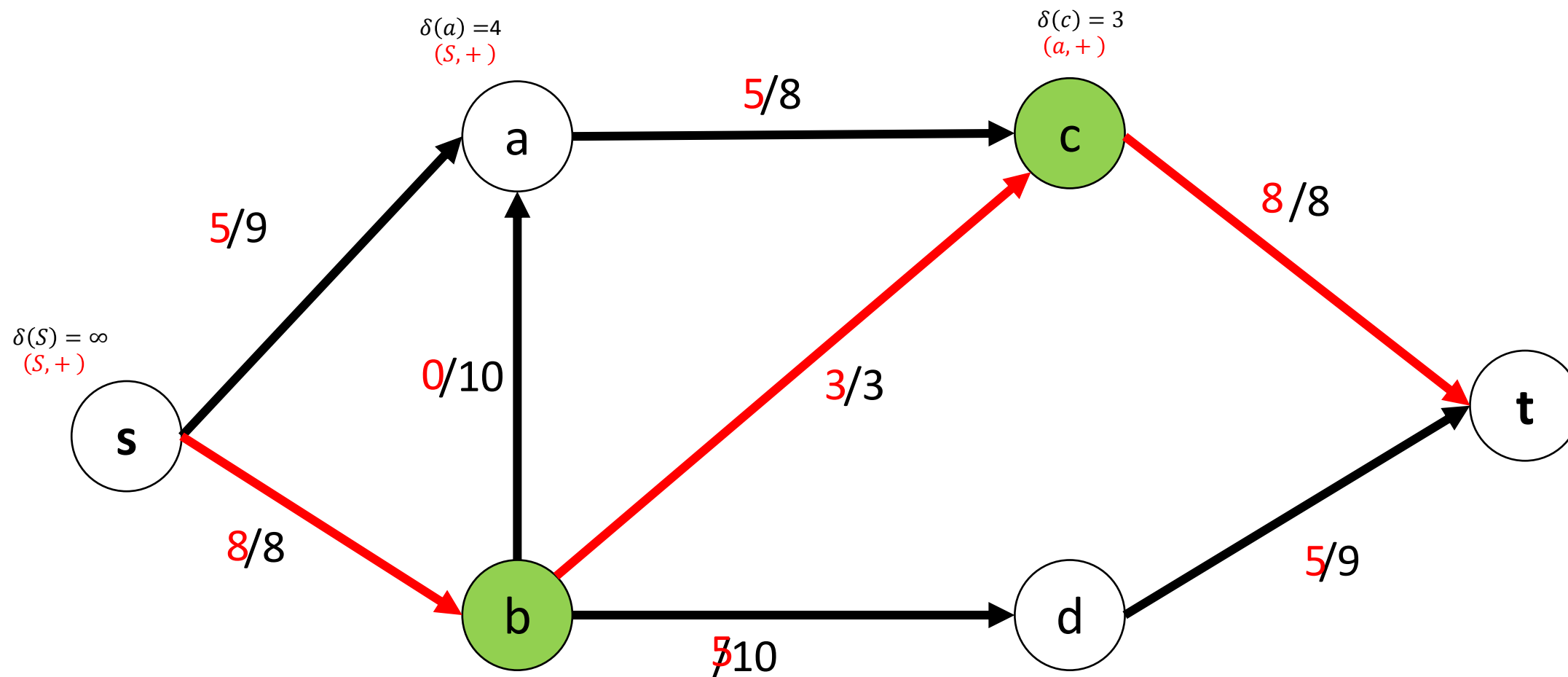


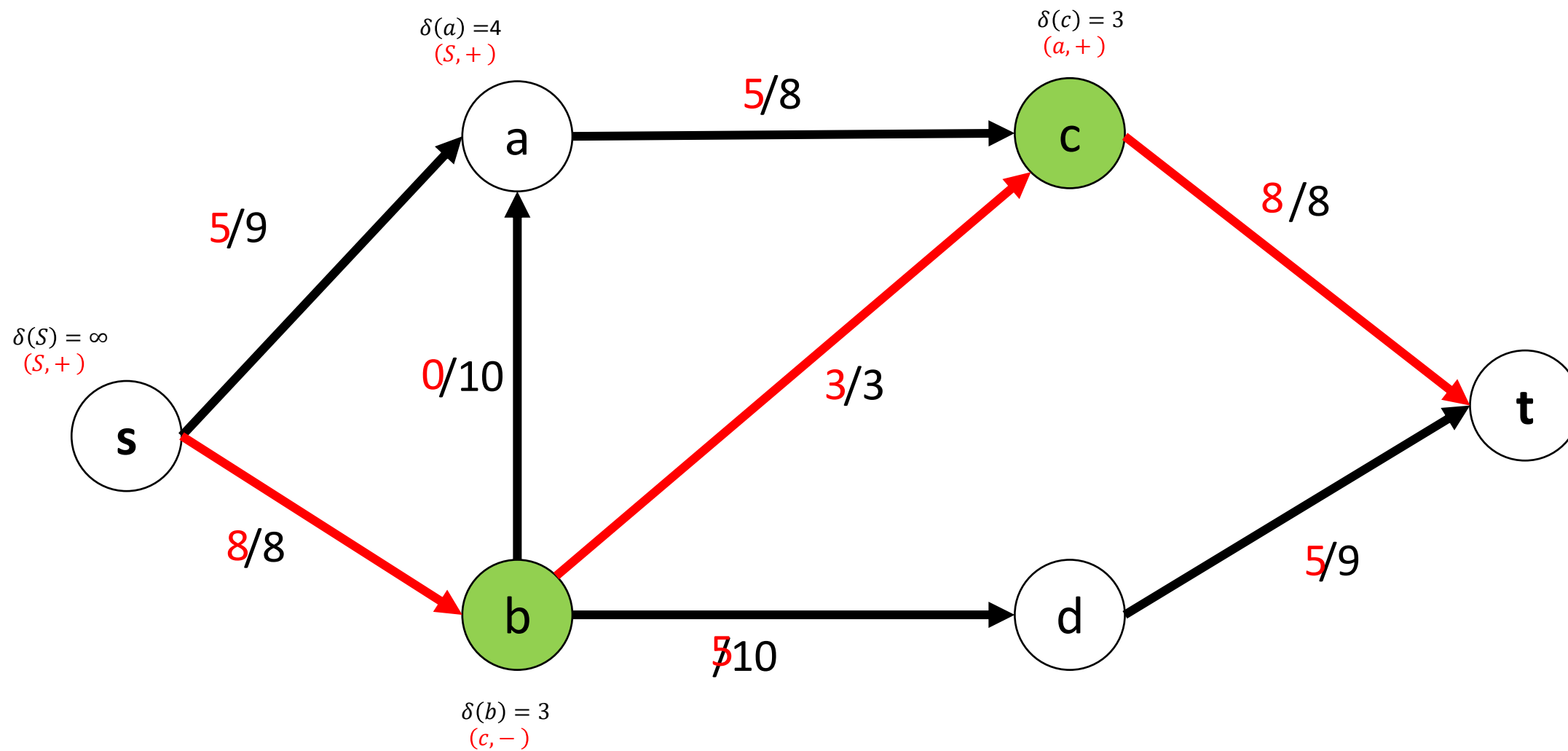


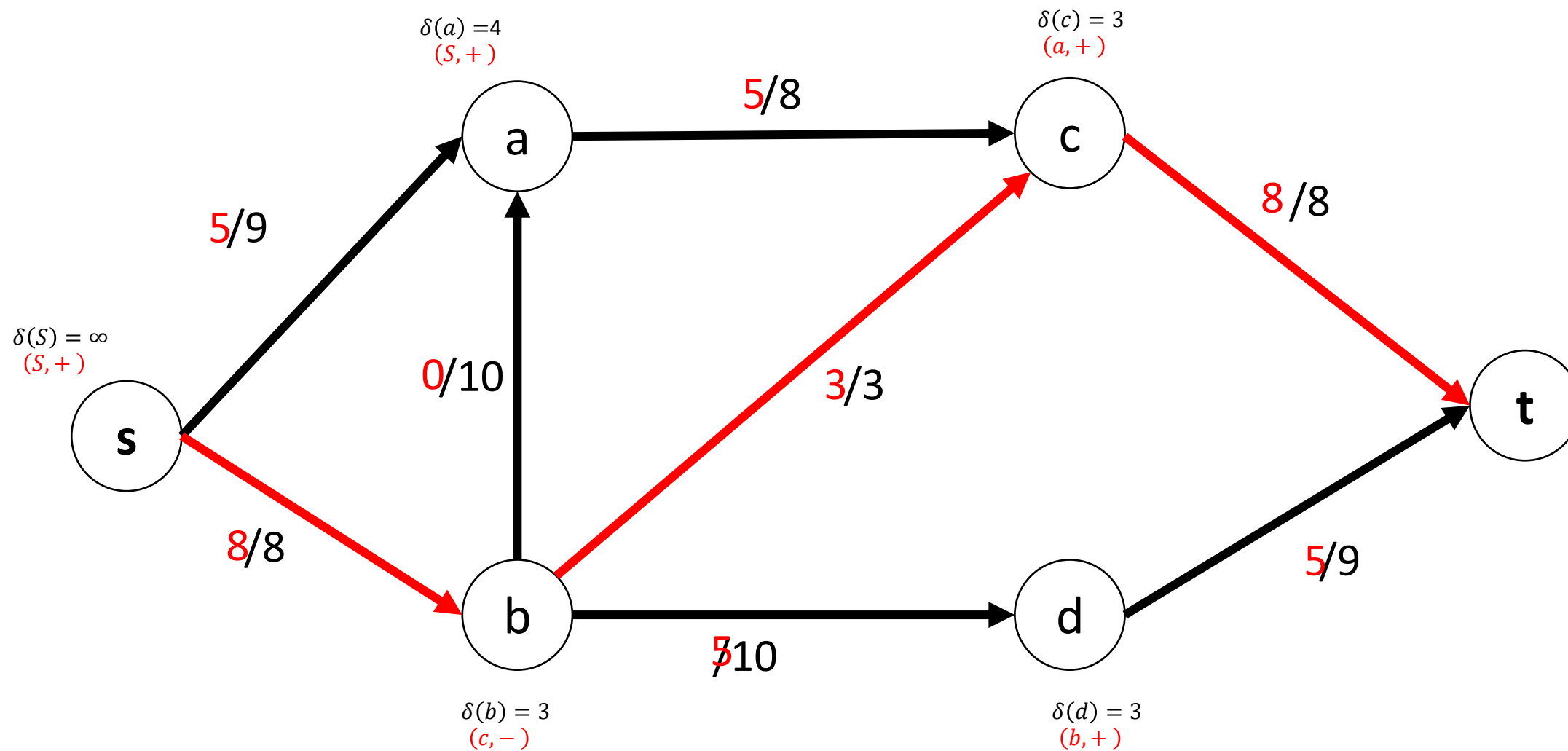


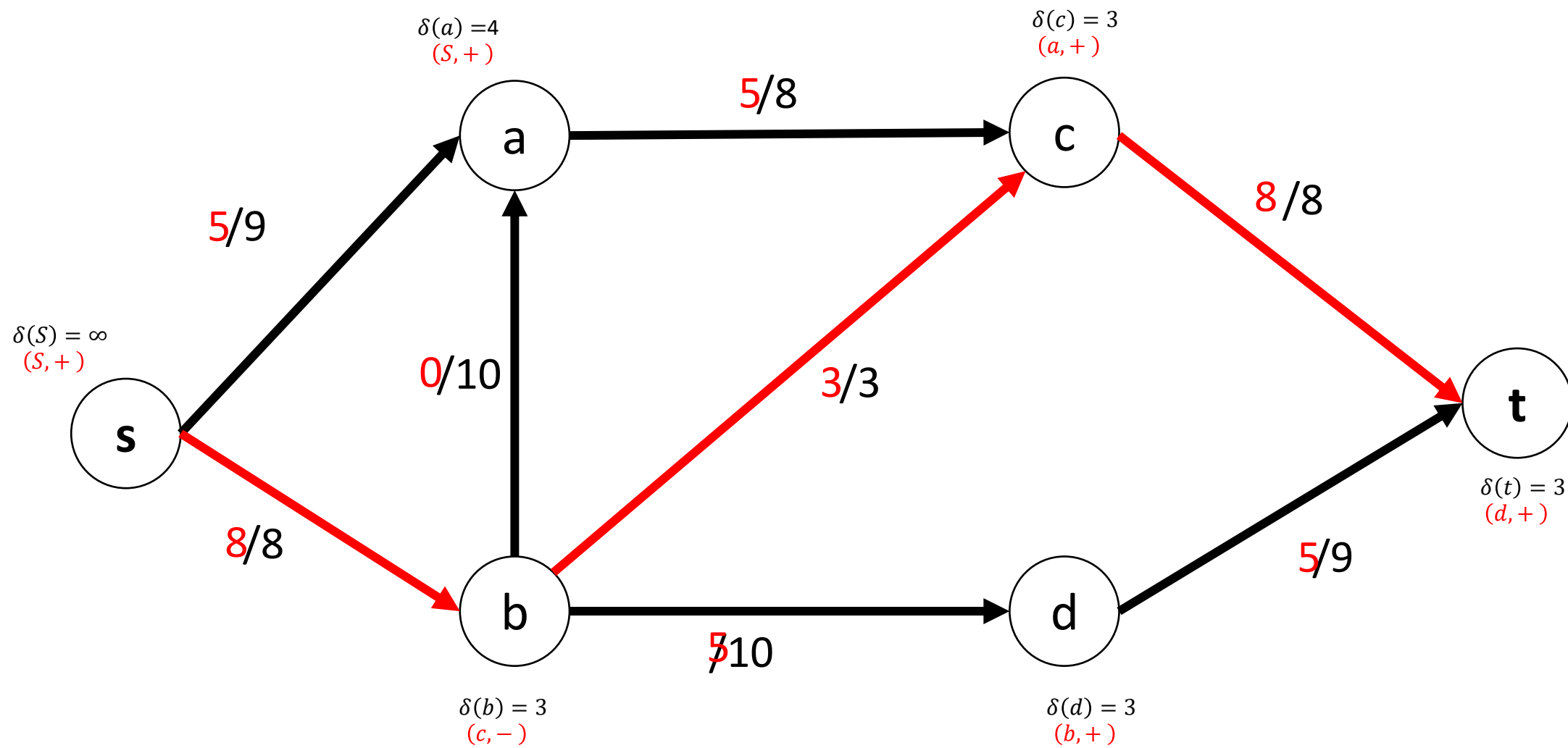


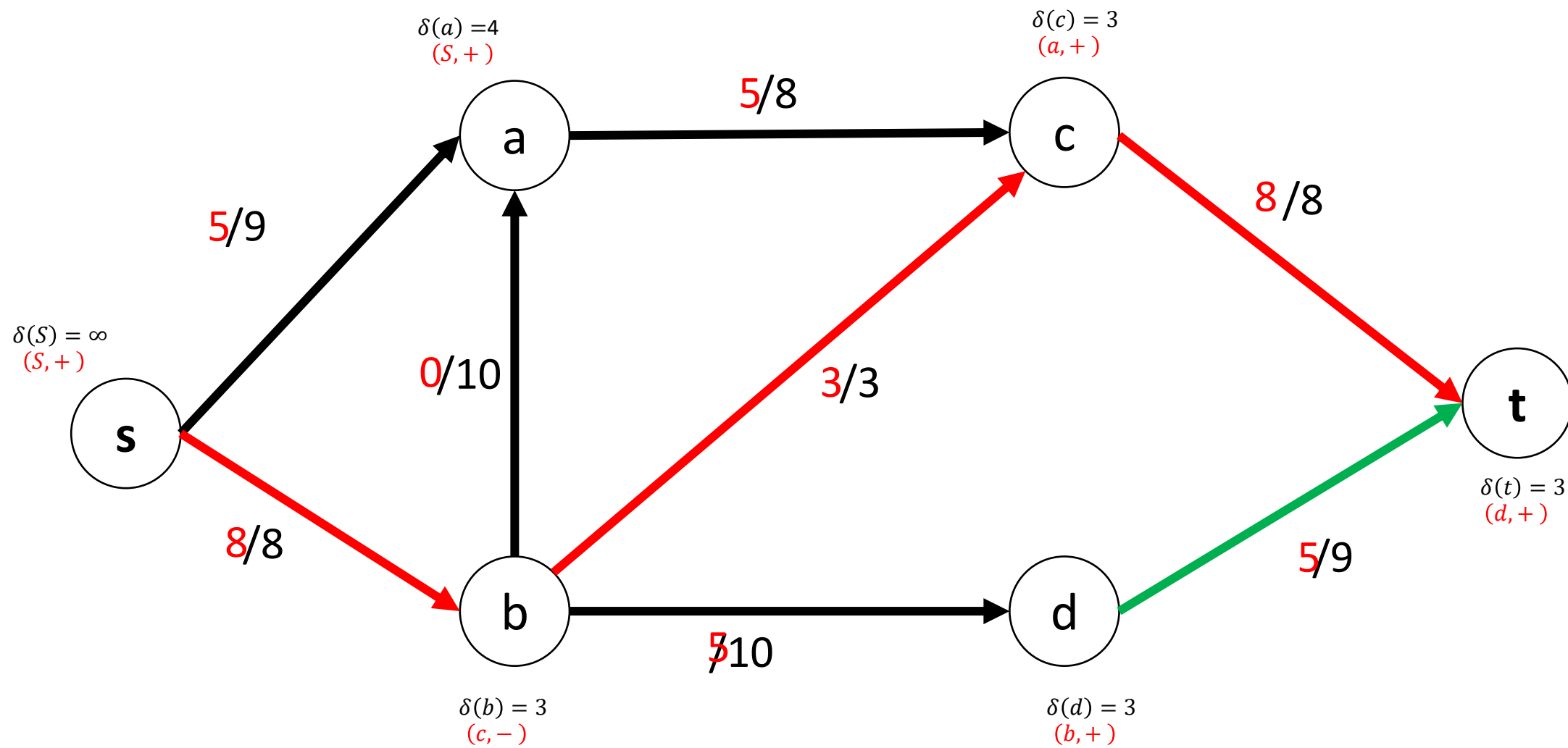


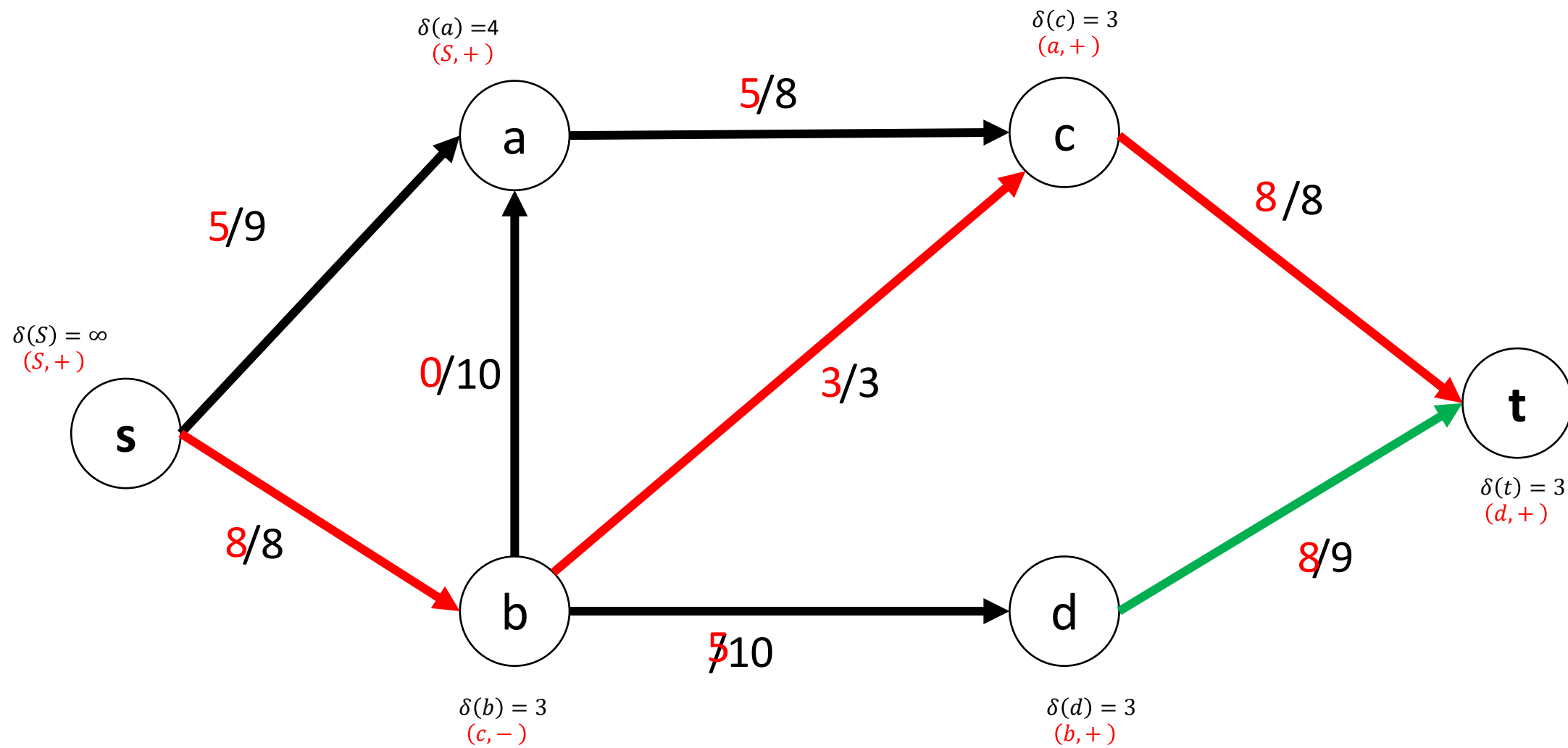


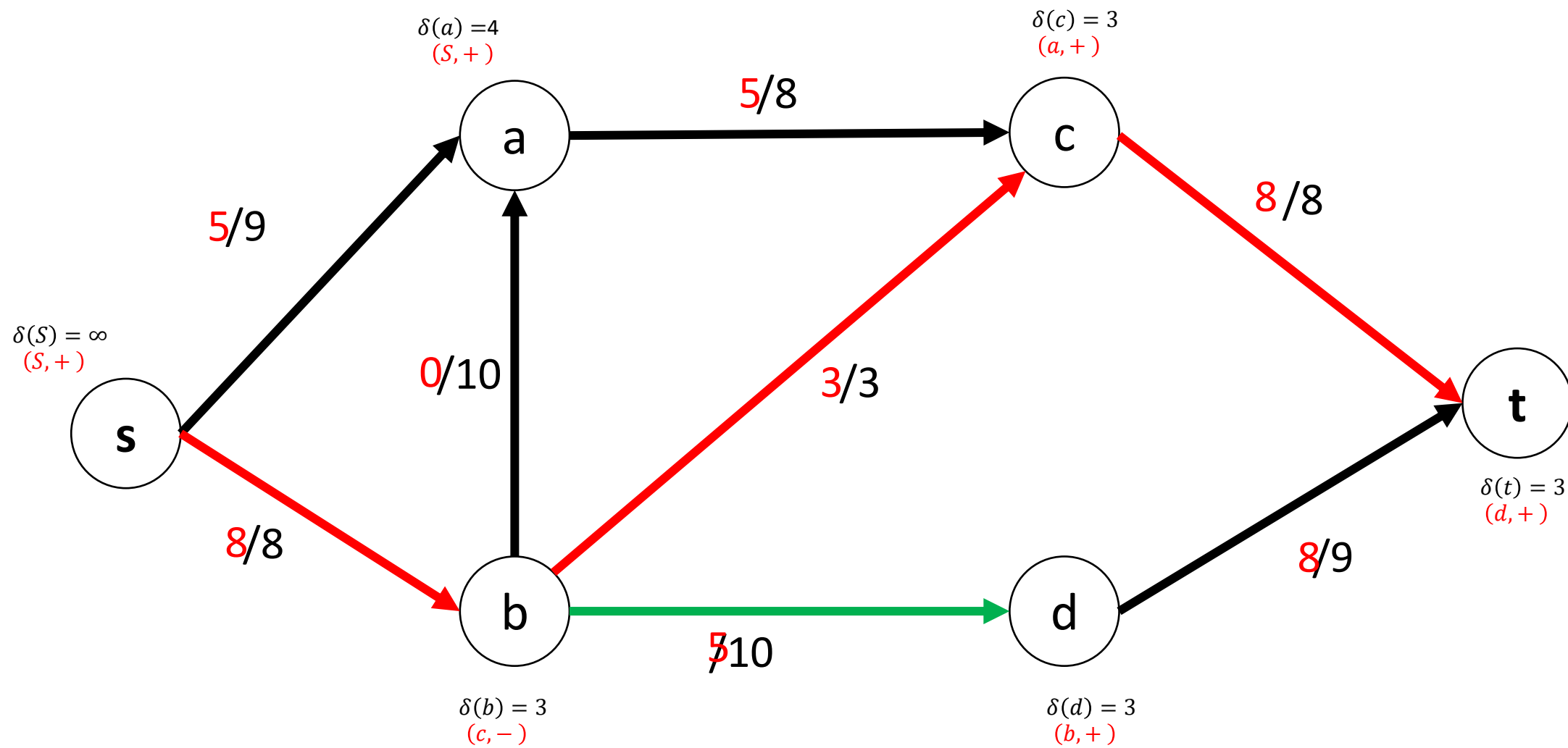


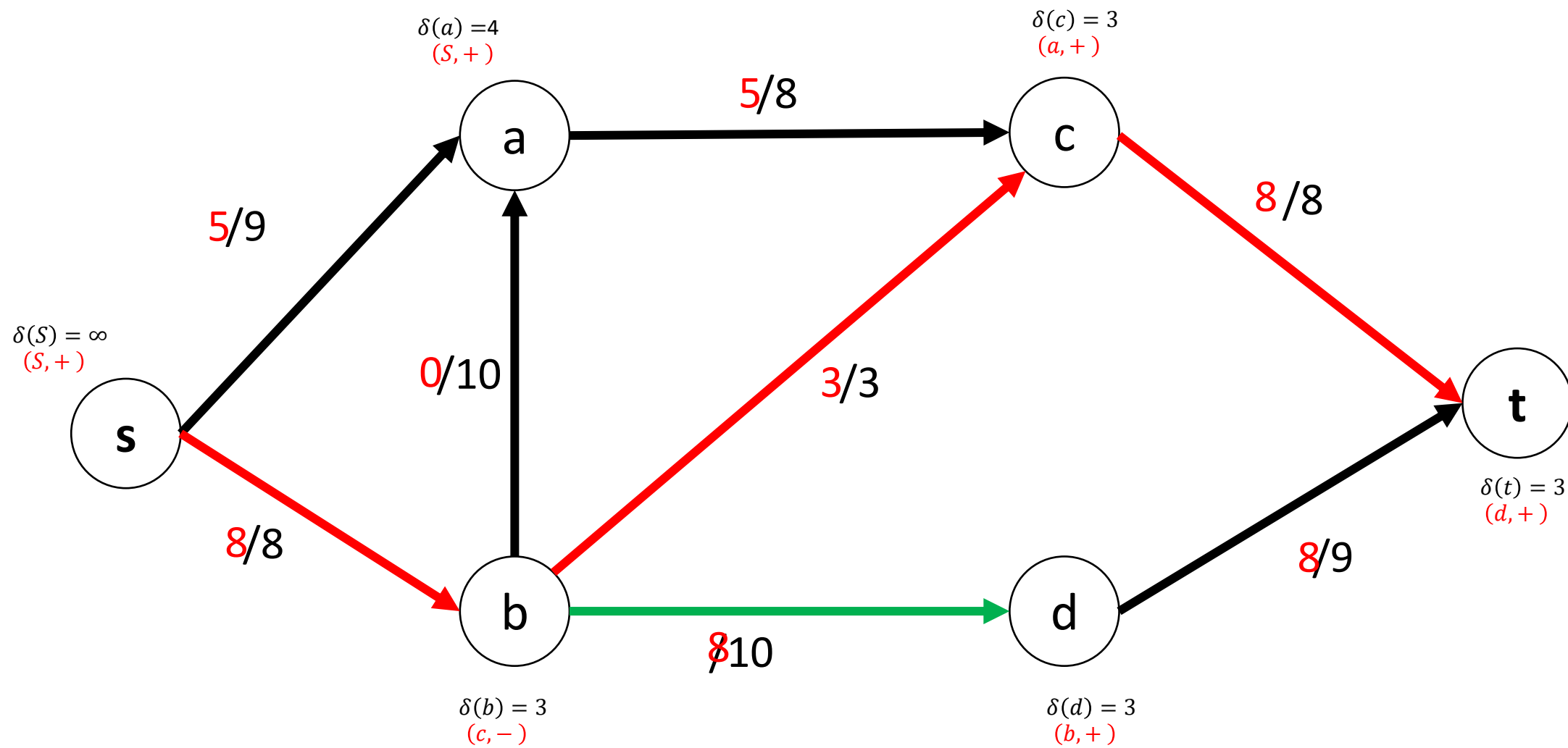


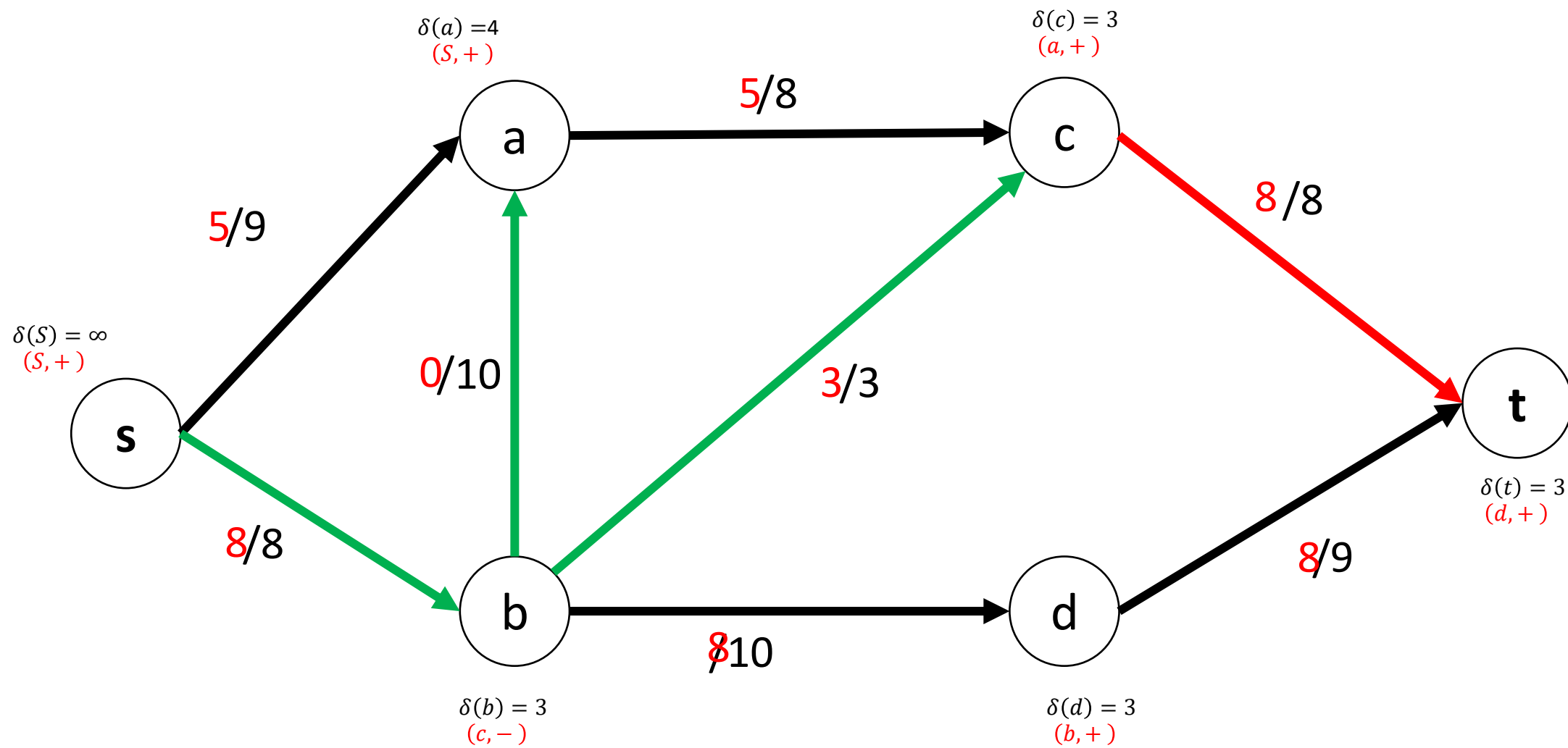


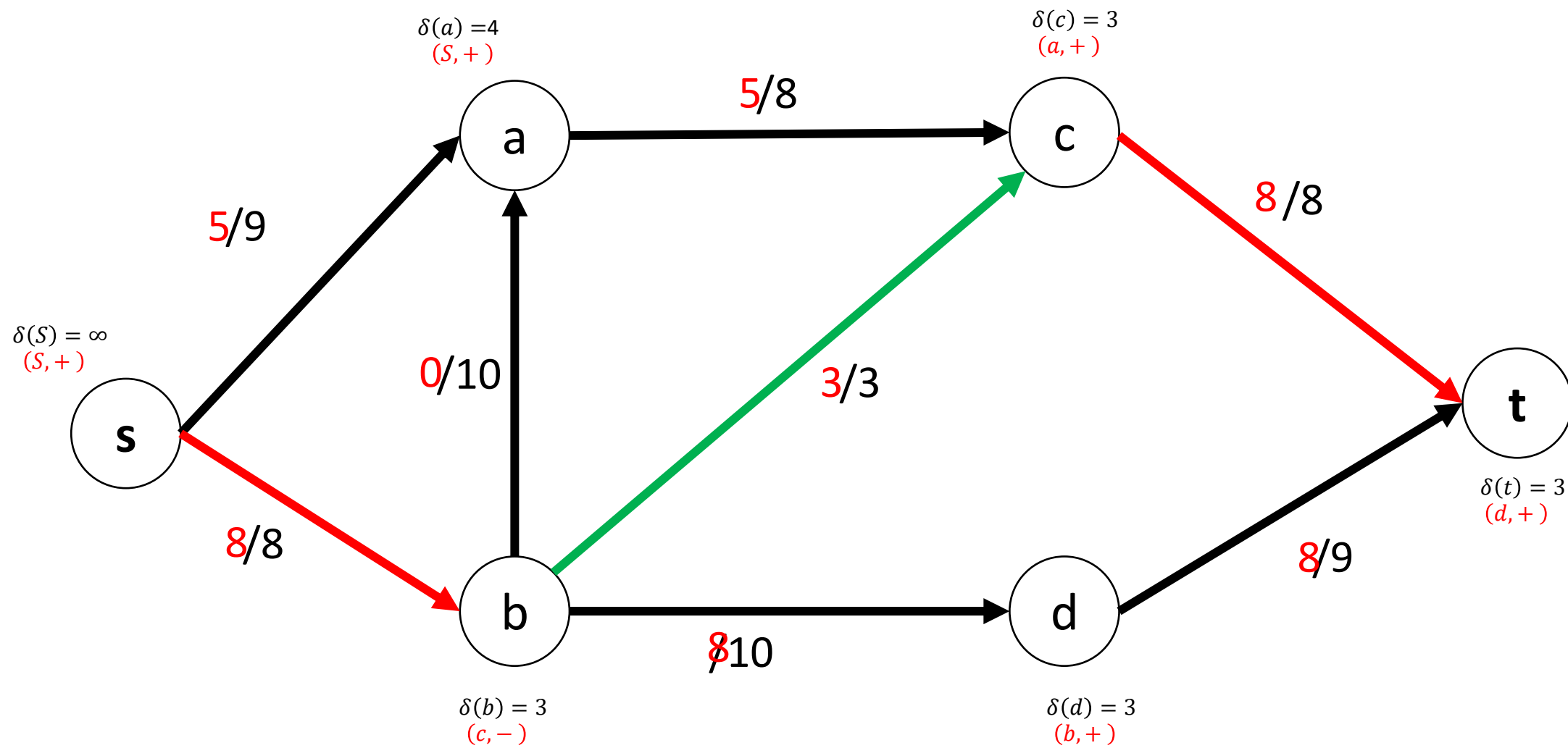


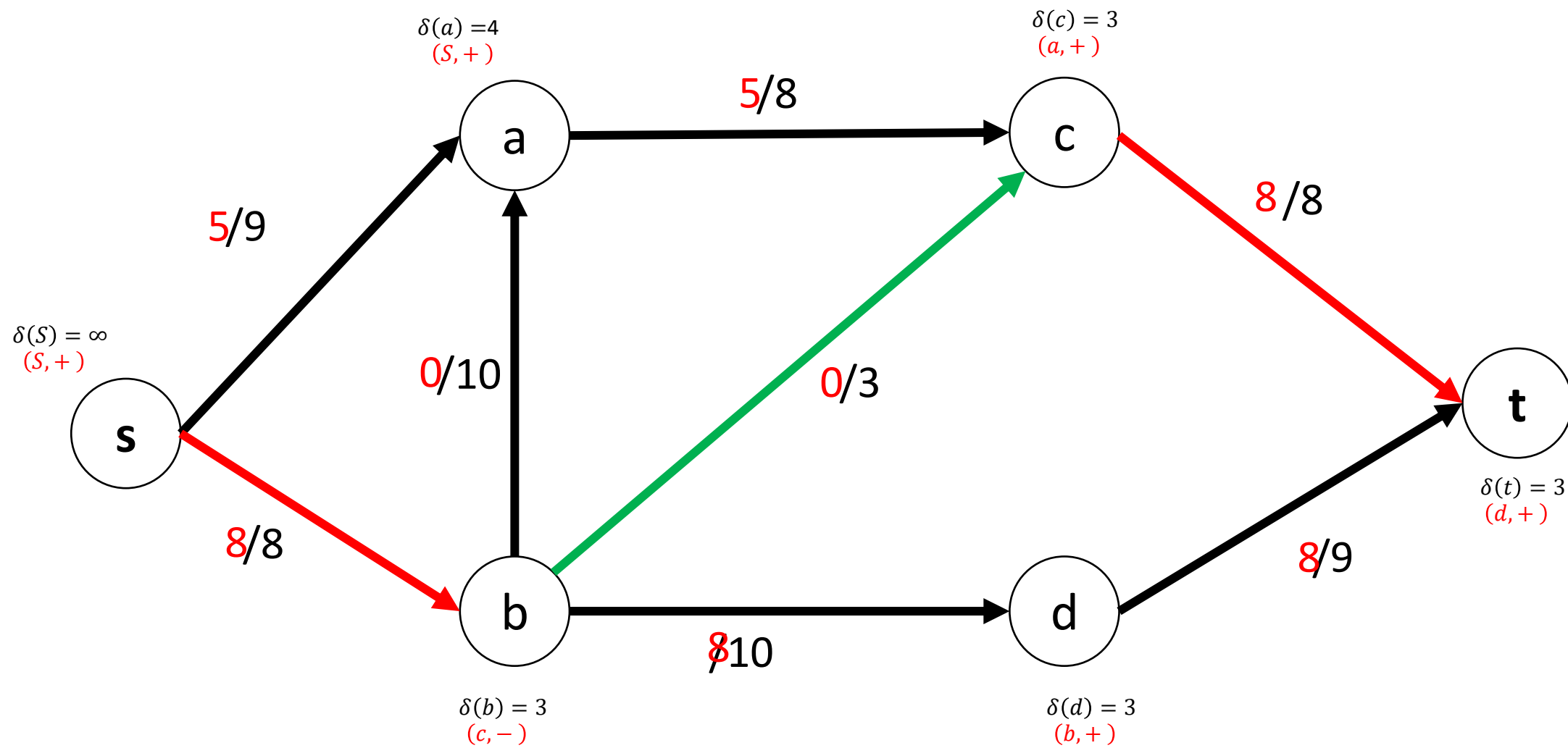


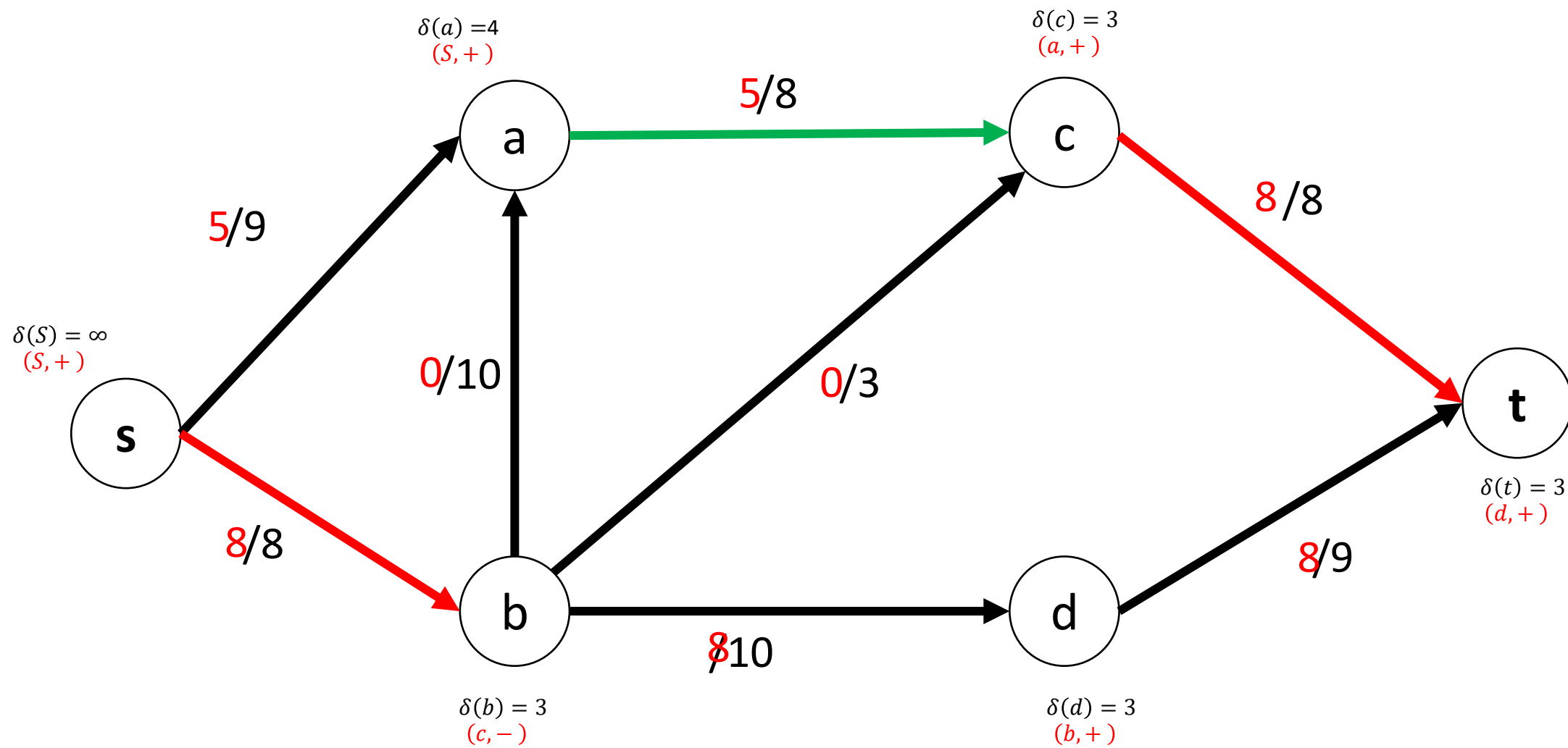


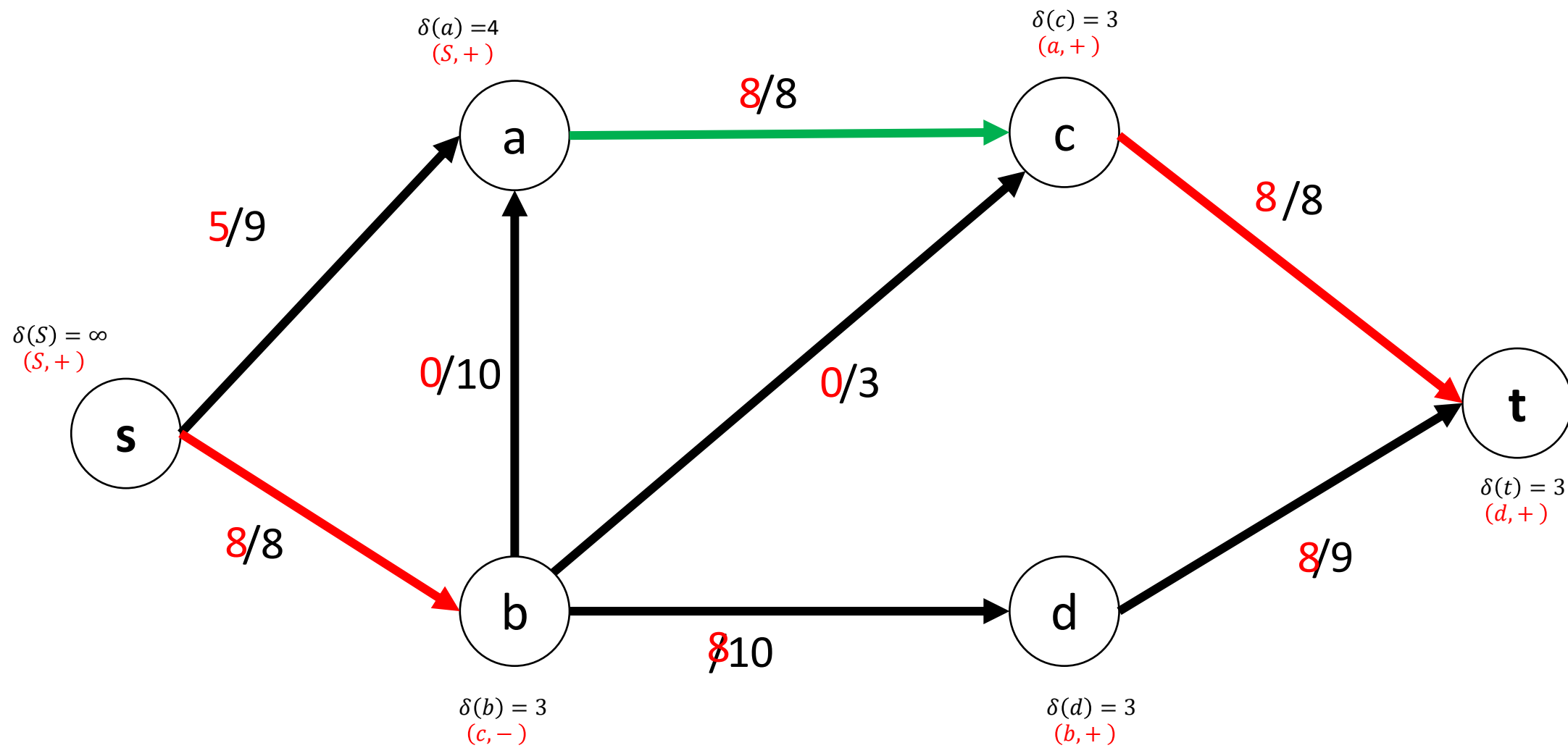


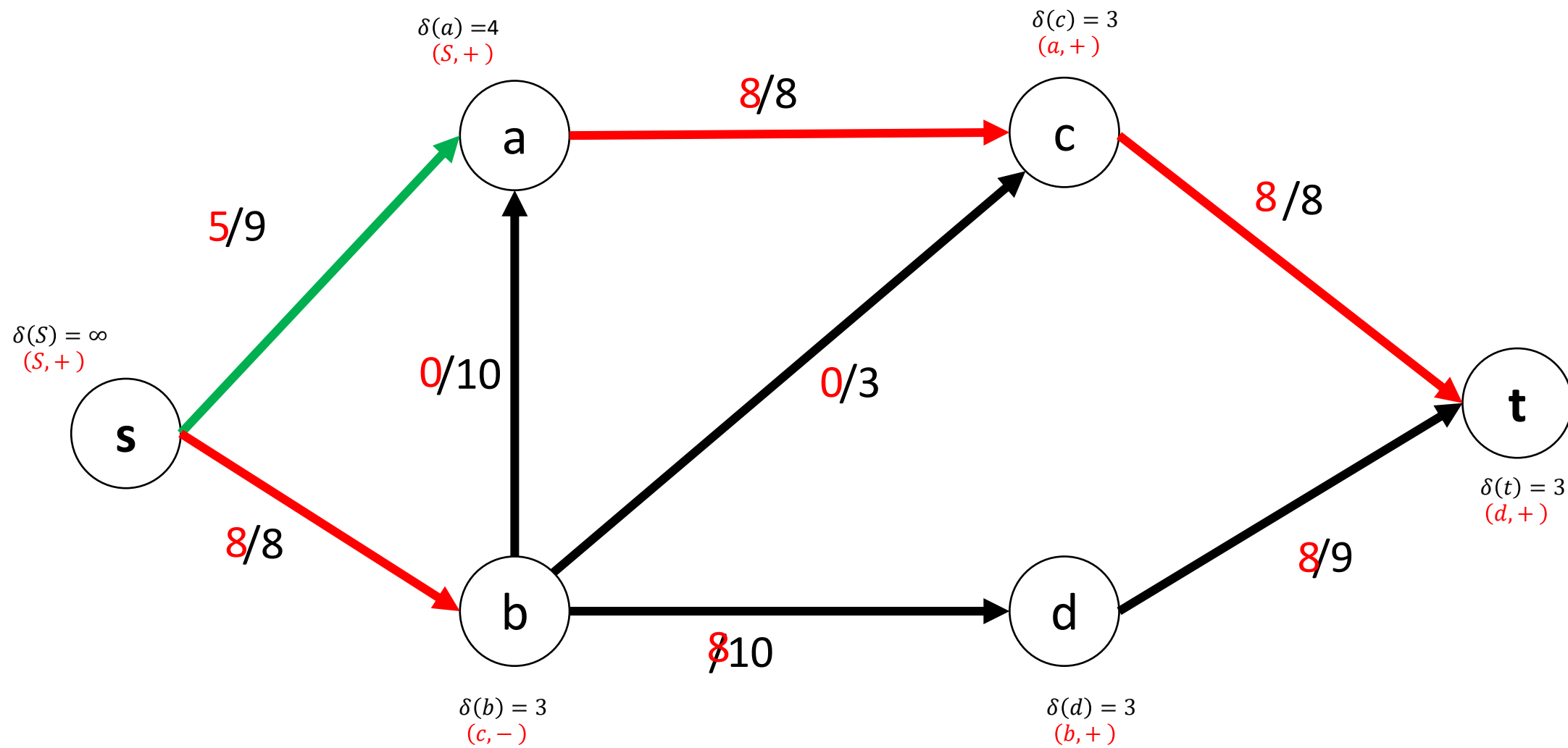


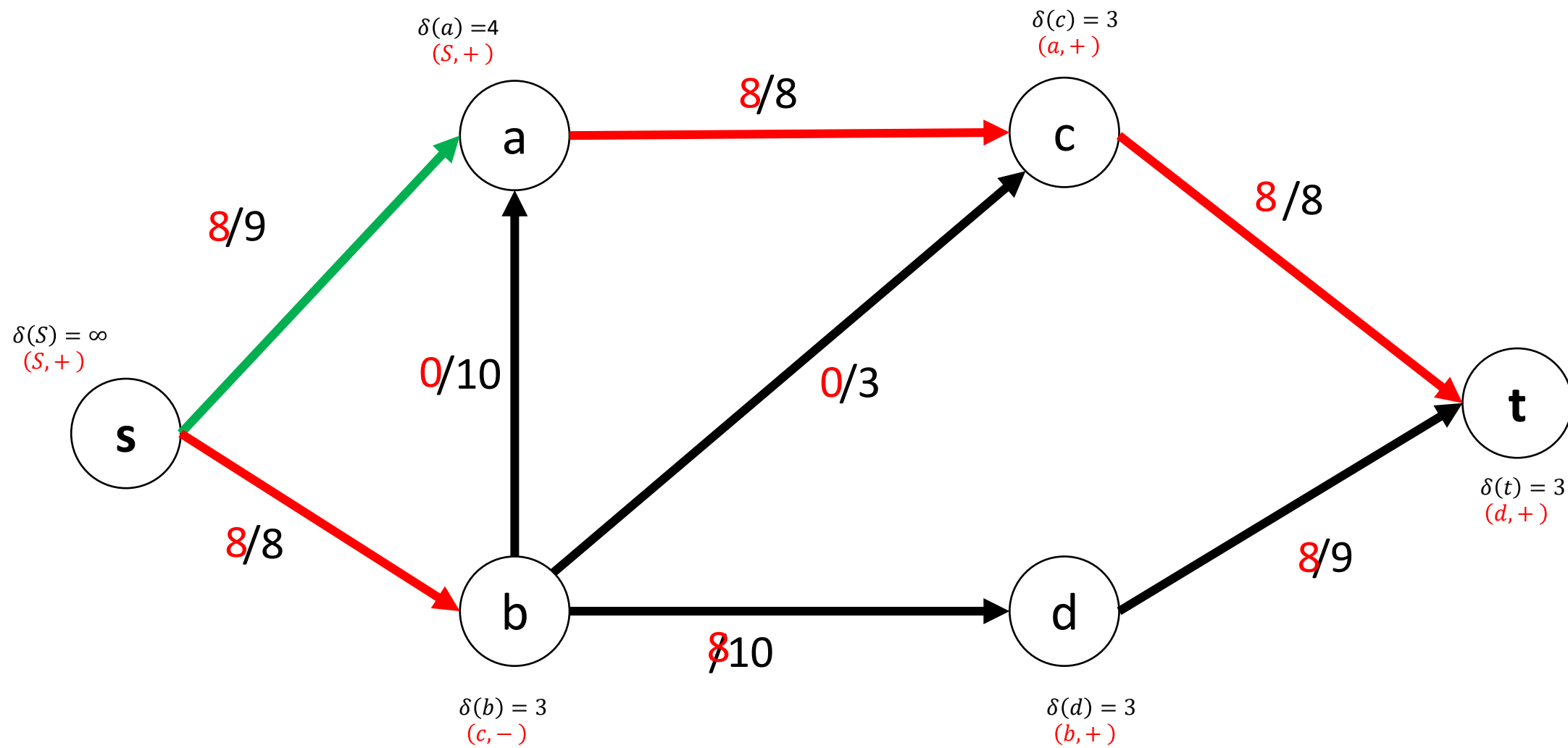


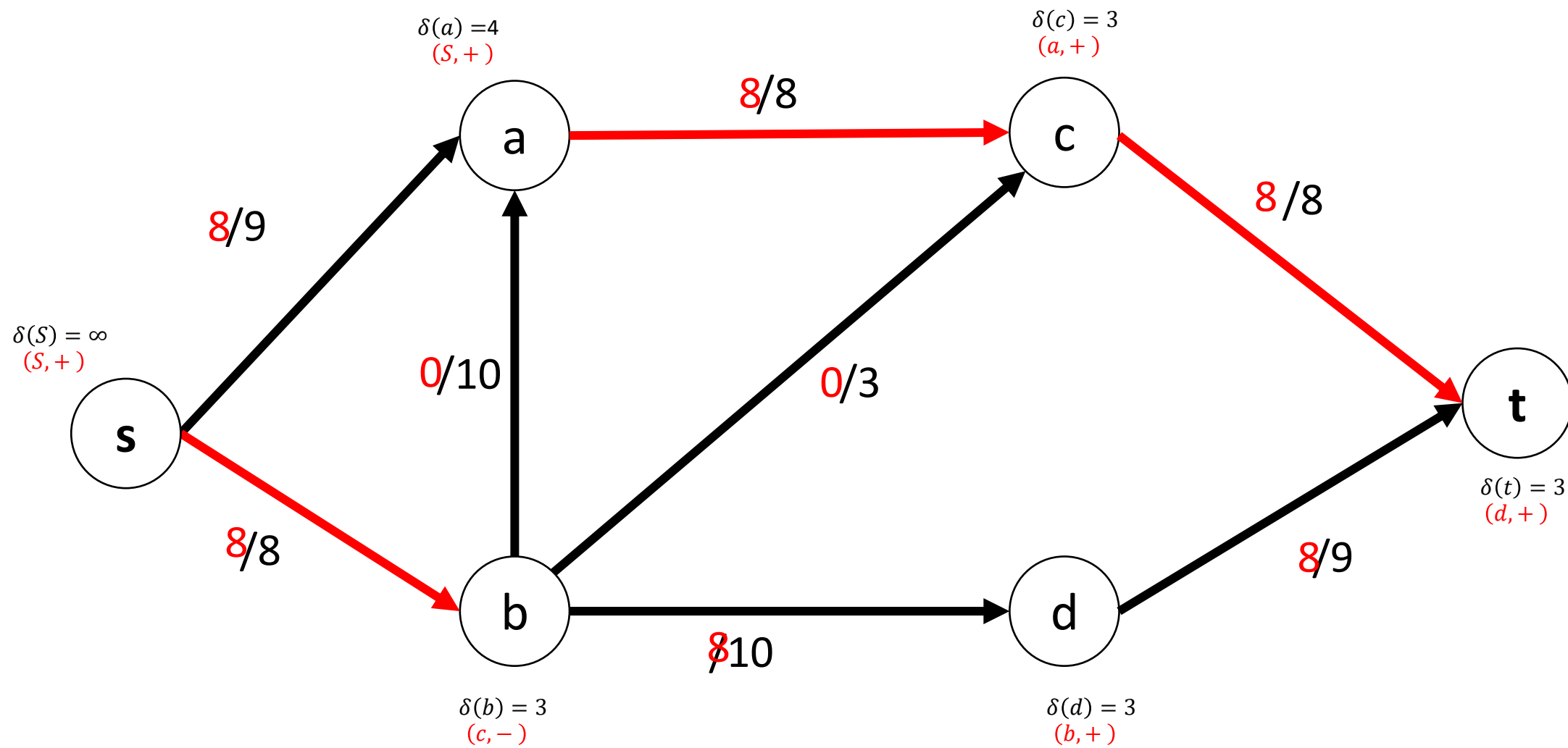


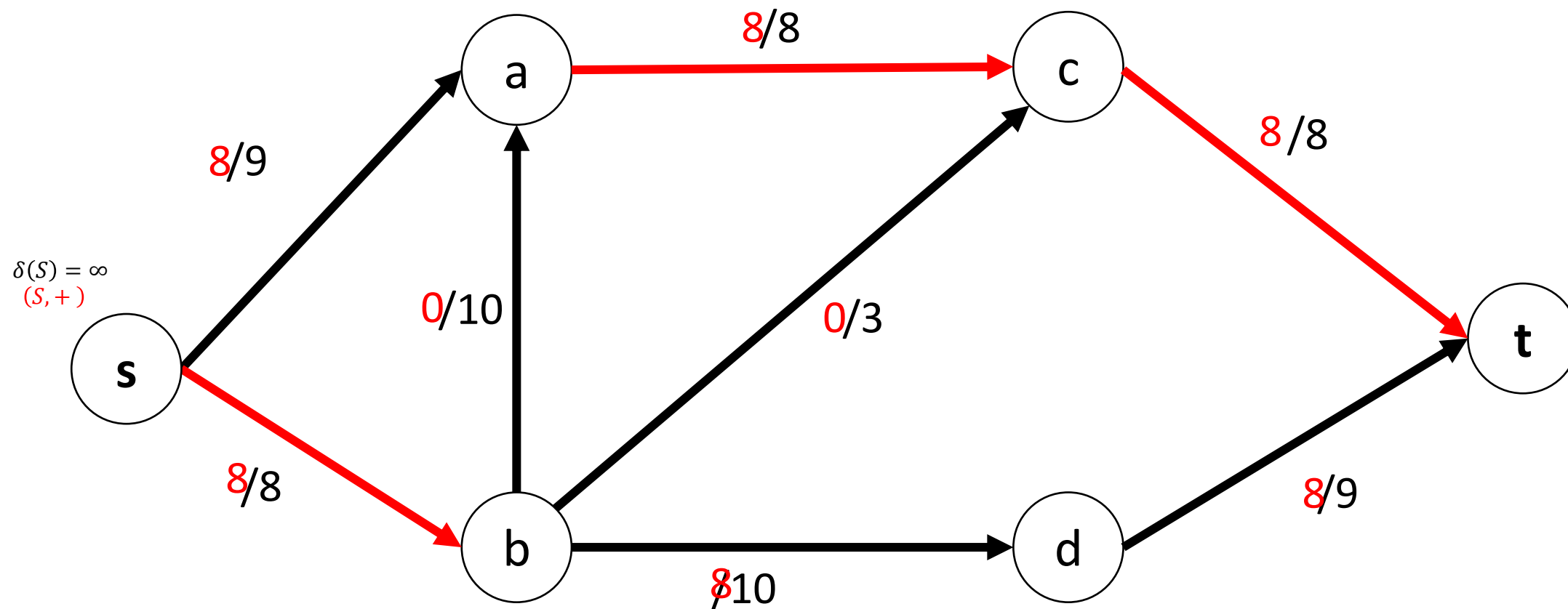


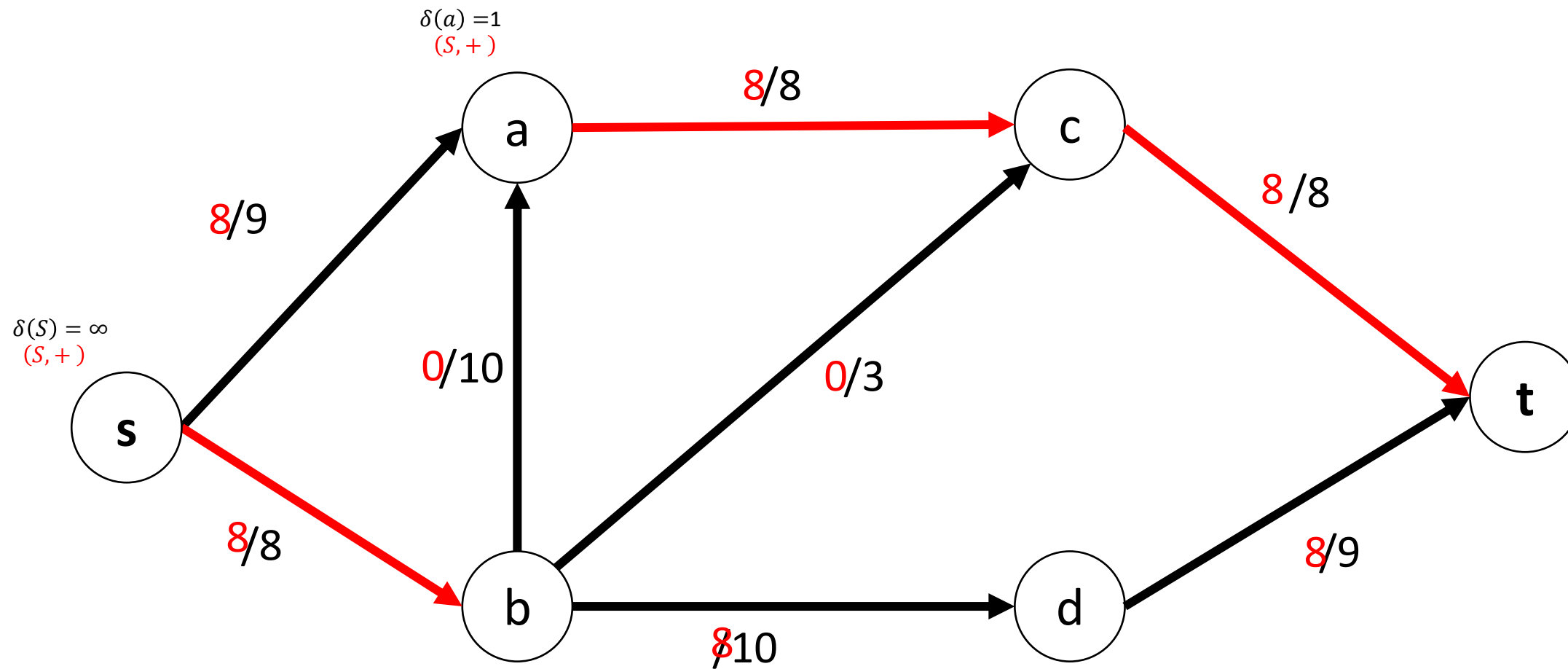


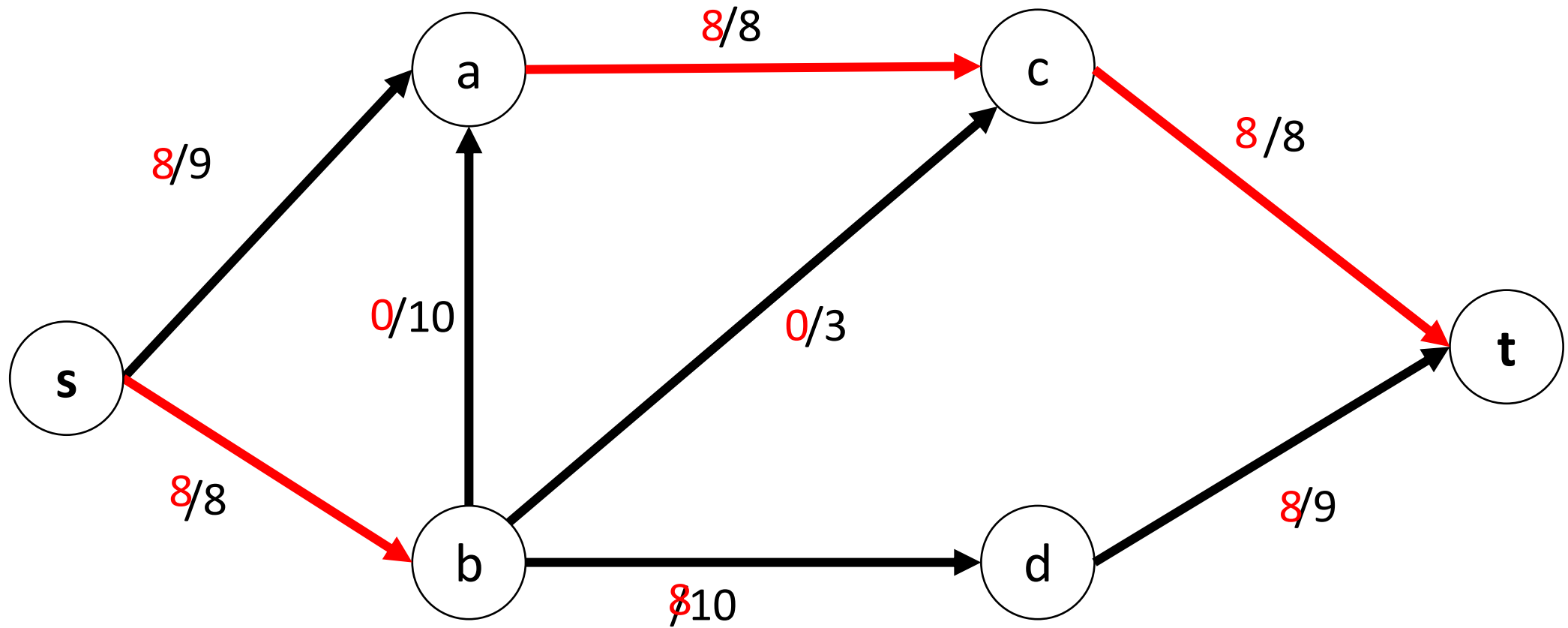




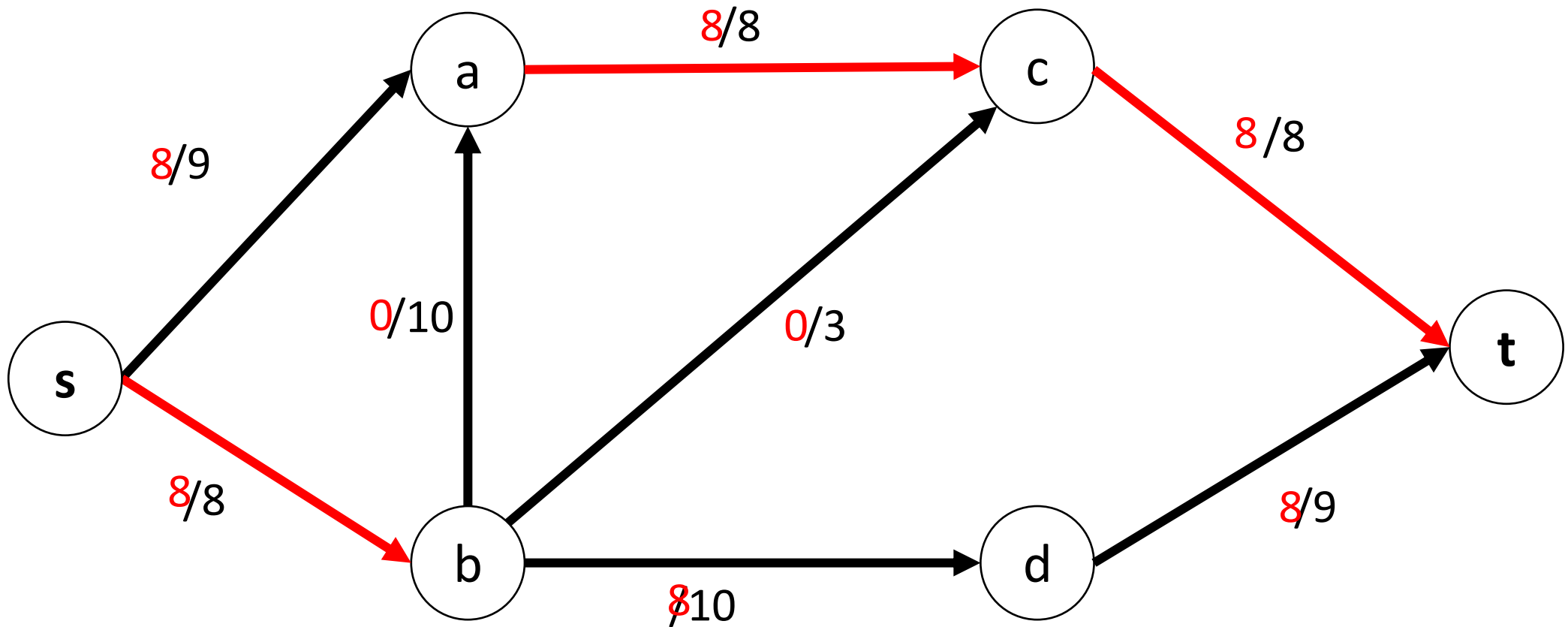








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$$S = \{s, a\}, S' = \{b, c, d, t\}$$

$$\delta(S) = \{(a, c), (s, b)\}$$

$$C(\delta(S)) = \sum_{(i,j) \in \delta(S)} C(i,j) = 8 + 8 = 16$$